

Introduction to Linear Algebra

5th Edition

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Summary

Six great theorem in LA:

Dimension Th. All bases for a vec sp. have the same number of vecs.

Counting Th. Dimension of row sp. + dimension of nullspace = numbers of cols.

Rank Th. Dimension of row sp. = dimension of col sp. This is the rank.

Fundamental Th. The row sp. and nullspace of A are ortho. complements in \mathbb{R}^n .

SVD There are ortho. bases (v 's and u 's for r and c sp.) so that $Av_i = \sigma_i u_i$.

Spectral Th. If $A^T = A$ there are ortho. q 's so that $Aq_i = \lambda_i q_i$ and $A = Q\Lambda Q^T$.

LA in a nutshell:

nonsingular

singular

A is inv. rows/cols are ind.

A is not inv. rows/cols are dependent.

$\det A$ is not zero.

$\det A$ is zero.

$Ax=0$ has one solution $x=0$.

$Ax=0$ has inf. many solutions.

$Ax=b$ has one solution $x=A^{-1}b$.

$Ax=b$ has no solution or inf. many.

A has n pivots. full rank $r=n$.

A has $r < n$ pivots. rank $r < n$.

rref $R=I$. row/col sp. is all of \mathbb{R}^n .

rref R has at least one zero row. row/col sp.

A 's all eigenvalues are nonzero.

zero is an eigenvalue of A . $\dim \leq r$.

$A^T A$ is sym. pos. has n (pos.) singular values.

$A^T A$ is only semipos. has $r < n$ singular values.

Matrix factorizations:

$$A=LU, A=LDU, PA=LU, EA=R, A=QR, A=XX^{-1}, S=Q\Lambda Q^T$$

$$(S=LDL^T)$$

$$A=U\Sigma V^T = (C^T C), A^T = V\Sigma^+ U^T, A=QS, A=U\Lambda U^{-1}, A=QTB^{-1}, A=QTQ^H$$

$$(U\Lambda U^H)$$

Alphabet

A any matrix

L lower tri. matrix

U upper tri. matrix, left singular matrix

E elimination matrix (, eye matrix)

P permutation matrix, pascal matrix, projection matrix

I identity matrix

D diagonal matrix

D_n (\bar{D}_n / \underline{D}_n) lower/upper difference matrix
(forward/backward)

R ref matrix, upper tri. matrix, reflection matrix, rotation matrix
(reduced row echelon form)

Q orthogonal matrix

C cofactor matrix (companion matrix), correlation matrix

S symmetric matrix

X eigenvector matrix

Λ eigenvalue matrix

V right singular matrix

Σ singular value matrix

J Jordan matrix

B basis matrix

1. Introduction to Vectors

1.1 vectors and linear combinations

lin. comb: $c\vec{v} + d\vec{w}$ fill a plane? solution: if right side is on the plane.

1.2 lengths and dot products

dot p: $\vec{v} \cdot \vec{w}$ perpendicular: $\vec{v} \cdot \vec{w} = 0$, length $\|\vec{v}\| = \sqrt{\vec{v} \cdot \vec{v}}$

all vectors have $|\vec{v} \cdot \vec{w}| \leq \|\vec{v}\| \|\vec{w}\|$ angle $\theta = \cos^{-1} \frac{\vec{v} \cdot \vec{w}}{\|\vec{v}\| \|\vec{w}\|}$

$\|\vec{e}_i\| = 1$ unit vector

1.3 matrices

$A\vec{x}$: comb of the cols, components of $A\vec{x}$ are the dots of A 's rows

equation: $A\vec{x} = \vec{b}$ solution could write $\vec{x} = A^{-1}\vec{b}$ and \vec{x} .

and some matrices don't have A^{-1} .

2. Solving Linear Equations

2.1 vectors and linear combinations $A\vec{x} = \vec{b}$

col picture: comb of n cols of A produces the vec \vec{b} .

row picture: m eqs from m rows give m planes meeting at \vec{x} .

$\vec{b} = 0$, at least one possibility is $\vec{x} = 0$.

2.2 the idea of elimination

$\begin{pmatrix} a_{11} \\ \vdots \\ a_{nn} \end{pmatrix}$ pivot a_{ii} , multiplier $l_{ij} = \frac{a_{ji}}{a_{ii}}$ and then subtract it from j .

elimination (Gauss ~) $A \rightarrow U$ elimination doesn't change solutions.

$A\vec{x} = \vec{b}$: forward elimination + back substitution

elimination breaks down if zero appears in pivot.

exchanging two eqs may solve it. (zero is not allowed in pivot \leftarrow nonsingular)

when breakdown is permanent, $A\vec{x} = \vec{b}$ has no solution or inf. many.

$A^2 = I \cdot A \cdot 0 \quad A^3 = I \cdot A \cdot 0$

幂等矩阵的性质:

(零/一) $A^2 = 0$: $C(A)$ is contained in $N(A)$
 so $r \leq n - r$ (秩)

$A^2 = A$: if idem. sym. mt, we call proj.
 (idempotent mt.) if idem. mt. full rank, it's inv.

$\text{tr}(A) = \text{rank}(A)$ $\oplus A \sim B, B \text{ idem.}$

$\det(A) = 0 \text{ or } 1. \quad A \sim \begin{pmatrix} I_r & \\ & 0 \end{pmatrix}$

$\lambda = 0, 1.$
 $(A^2 = A \text{ odd idem. } \lambda = 0, \pm 1)$
 \downarrow vice versa!



引 \times LA?

1. 线性空间 $\vec{x} = \alpha\vec{u} + \beta\vec{v} + \gamma\vec{w}$

$\vec{x} = \begin{pmatrix} \alpha \\ \beta \\ \gamma \end{pmatrix} \quad \vec{u} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$ 向量

基底. 线性变换与变基

$A\vec{x} = \vec{b}$

inv. 变换

$\lambda = 0, 1.$

2. 分块初等 $\begin{pmatrix} A & \\ & B \end{pmatrix}$ 多向量 AB

矩阵乘法

\downarrow 对应

$(A^2 = A \text{ odd idem. } \lambda = 0, \pm 1)$
 \downarrow vice versa!

求行列式

直算/余子式/降阶 Big Formular Δ 递推 Δ 全排列

消元 \rightarrow 上下三角

大分拆 Big Partition Δ 剪子阵 $\begin{pmatrix} \times & & \\ & \times & \\ & & \times \end{pmatrix} = \begin{pmatrix} \times & & \\ & \times & \\ & & \times \end{pmatrix}$

加边



shear matrix

特殊阵

Vandermonde $V_n = \prod_{i>j} (\alpha_i - \alpha_j)$

cycle (冲顶 even)

tridiagonal (band) triD det: $D_n = aD_{n-1} - bcD_{n-2}$

(见 more properties.)

求逆

直算 A^{-1}

解方程 $AA^{-1} = I$

初等变换/加边 $E \leftrightarrow E^{-1}(L)$

(消元)

$(A|I) \leftrightarrow (I|A^{-1})$ — G-J elimination

分块

$\begin{pmatrix} A & \\ & B \end{pmatrix}$

Δ Sherman-Morrison 公式: $(A + uv^T)^{-1} = A^{-1} - \frac{A^{-1}uv^T A^{-1}}{1 + v^T A^{-1}u}$

(见 2.4)

等式变形 $A^n = I \cdot f(A)$

esp. $(I + uv^T)^{-1} = I - \frac{uv^T}{1 + v^T u}$

左行右列

行 \vec{I}_n 列仿证

矩阵等式 (张量代数) 证明

非方阵不同时存在左右逆: $A \text{ } m \times n$

(广义逆)

left inv., right inv. $n \times m$

comb 出 $\begin{pmatrix} \vdots \\ \vdots \end{pmatrix} \dots \begin{pmatrix} \vdots \\ \vdots \end{pmatrix} \Rightarrow \text{rank } A \geq n, \text{rank } A \geq m$

而 $\text{rank } A \leq n, m$ 故 $n = m$.

左逆 = 右逆

left $\begin{cases} A\vec{x} = \vec{b} \Rightarrow \vec{x} = A_l^{-1}\vec{b} \text{ 唯一解} \\ A\vec{x} = (AA_l^{-1})\vec{b} = \vec{b} \\ \forall \vec{b} \Rightarrow AA_l^{-1} = I, A_l^{-1} = A_r^{-1} \end{cases}$

right $\begin{cases} \vec{x}_0 = A_r^{-1}\vec{b} \text{ 为一解} \\ A_r^{-1}A\vec{x}_0 = \vec{x}_0 \quad \forall \vec{b} \Rightarrow \forall \vec{x}_0 \\ A_r^{-1}A = I, A_r^{-1} = A_l^{-1} \end{cases}$

inv. uniqueness: $A_l^{-1} = A_r^{-1} \cdot I = A_l^{-1}AA_2^{-1} = A_2^{-1}$

如果非 \vec{v} ,
 则无法有
 $A\vec{x} = \vec{b}$

2.3 elimination using matrices

views of matrix multiplications: $A = (\bar{a}_1, \dots, \bar{a}_n)$, $A\bar{x} = x_1\bar{a}_1 + \dots + x_n\bar{a}_n$

$E(A\bar{b}) = (EA, E\bar{b})$ augmented matrix $AB = (A\bar{b}_1, \dots, A\bar{b}_n)$

elimination matrix / elementary transform (ET) : $E_{ij} = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & -1 \end{pmatrix}$

permutation matrix $P_{ij} = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & 1 \end{pmatrix}$ (may used) \rightarrow could see from ① multiply / act vecs. rows/cols

one matrix both steps: $\prod P E \dots = E$
times then. to create a $U = EA$

2.4 rules for matrix operations

$(m \times n)(n \times p) = (m \times p)$ mnp separated muls

$(AB)C = A(BC)$. $AB \neq BA$ usually

block multiplication, if their shapes permit.

block elimination. Schur complement:

$\begin{pmatrix} A & B \\ C & D \end{pmatrix} \rightarrow \begin{pmatrix} A & -A^{-1}B \\ 0 & D - CA^{-1}B \end{pmatrix}$
so Ax \rightarrow recomb of A cols \rightarrow x after a linear transform (bases (dual) transform)

- \rightarrow A rows \cdot B cols (mp entries)
- \rightarrow A cols \cdot B rows (n adds)
- \rightarrow A rows \cdot B (vec. mt)
- \rightarrow A \cdot B cols \star (mt. vec)
- \rightarrow B rows recomb (comb)
- \rightarrow A cols recomb (comb)
- \rightarrow a left row, right col!

2.5 inverse matrices

invertible: $AA^{-1} = A^{-1}A = I$. (means $n \times n$) test invertibility:

(not sq. mt do not have inv., for transpass Δ log:

if all invertible, $(AB)^{-1} = B^{-1}A^{-1}$

Gauss-Jordan elimination: $O_{(G-J)} = I_n$

$(AI) \rightarrow (IA^{-1})$ $(A\bar{b}) \rightarrow (I\bar{x})$

(Jor means back elimination, the same essence)

$\square \begin{pmatrix} 1 & 0 \\ -5 & 1 \\ 0 & -1 \end{pmatrix}^{-1} = \begin{pmatrix} 1 & 0 \\ 5 & 1 \\ 0 & -1 \end{pmatrix}$

so: how you view a trint.?

- ① A must have n (nonzero) pivots.
- ② $\det A$ must not be zero.
- ③ $A\bar{x} = 0, \bar{x} = 0$ must be the only solution.
- ④ $A^{-1}A\bar{x} = \bar{b}, \bar{x} = A^{-1}\bar{b}$ be the only solution, not none or inf.)
- ⑤ rows/cols of A are independent.
- ⑥ row/col space are \mathbb{R}^n . $\oplus^n \text{rank } A = n$

$\circ \bar{x} = A^{-1}\bar{b}$'s a math trans!

\star elim⁻¹ is no interception (coupling, rolling)
elim is rolling.

1. acts on row or col? / how it produced?
2. view as E or directly E^{-1} ?
3. start from tip or top? one can be rolling.

to get an inv., must go through rolling once.
(turn on unrolling way back)

矩阵相乘 ($A=B \cdot inv.$) 与行列空间的基 (comb)

blockize: better
sq. in diag.

• useful justify dependent/uninv.: r/c s adds to zero, or one has combs of others.

• 矩阵收集/包:

• adjacency matrix $S = \begin{pmatrix} 0 & & \\ i & 0 & \\ & & \ddots \end{pmatrix}$

(无向图为 sym.) $S^2 \dots S^n$: length n , node $i \rightarrow j$

• sum matrix $S = \begin{pmatrix} & & \\ & & \\ 1 & & 0 \end{pmatrix}$

• difference matrix $\bar{D} = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & 0 \end{pmatrix}$ $D = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & 0 \end{pmatrix}$ $D^{-1} = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & 0 \end{pmatrix}$ or $\begin{pmatrix} & & \\ & & \\ 0 & & 1 \end{pmatrix}$

n .th diff matrix D_n

$$D_n^{-1} = S^n$$

• cycle matrix (zero-cycle matrix means adds equal to zero,

$C = \begin{pmatrix} a & b & \dots \\ b & c & \dots \\ \dots & \dots & \dots \end{pmatrix}$ it's uninv.)

• Pascal matrix / binomial coeff. matrix

$$P_L = \begin{pmatrix} 1 & & & \\ 1 & 1 & & \\ 1 & 2 & 1 & \\ 1 & 3 & 3 & 1 \end{pmatrix}$$

$$D = \begin{pmatrix} 1 & & & \\ & -1 & & \\ & & -1 & \\ & & & -1 \end{pmatrix}$$

$$P_L D P_L D = I$$

$$D P_L D = \begin{pmatrix} 1 & & & \\ -1 & 1 & & \\ -1 & -2 & 1 & \\ -1 & -3 & -3 & 1 \\ 1 & -4 & -6 & -4 & 1 \end{pmatrix}$$

sym. pascal matrix $P_{sym} = \begin{pmatrix} 1 & 1 & 1 & 1 & \dots \\ 1 & 2 & 3 & 4 & \dots \\ 1 & 3 & 6 & \dots & \dots \\ 1 & 4 & \dots & \dots & \dots \end{pmatrix} = P_L P_U$ (so $|P_{sym}| = 1$)

• diagonally dominant matrix $a_{ii} > \sum_{j \neq i} a_{ij}$ is inv.

for $A\vec{x} = 0$ only can be zero.

□ cyclic matrix

循环矩阵 (环)

$$C = \begin{pmatrix} a_1 & a_2 & a_3 & \dots & a_n \\ a_n & a_1 & & \dots & a_{n-1} \\ a_{n-1} & a_n & & \dots & \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_2 & a_3 & \dots & \dots & a_1 \end{pmatrix}$$

• basic cyclic matrix

$$J = \begin{pmatrix} 0 & 1 & & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 1 & 0 & \dots & 0 \end{pmatrix}, J^k = \begin{pmatrix} 0 & & & I_{n-k} \\ & \ddots & & \\ & & I_k & \\ & & & 0 \end{pmatrix} \quad (k \bmod n)$$

$$C = a_1 I_n + a_2 J + \dots + a_n J^{n-1} = g(C) \text{ eigenfunction.}$$

• eigen value

$$|\lambda I_n - J| = \lambda^n - 1$$

$\{I_n, J, \dots, J^{n-1}\}$ a base in \mathbb{C}^n

$1 \leq k \leq n-1$ (univ. root) $\omega_k = \cos \frac{2k\pi}{n} + i \sin \frac{2k\pi}{n}$ eigenvector $\alpha_k = (1, \omega_k, \omega_k^2, \dots, \omega_k^{n-1})^T$

$$X = \begin{pmatrix} 1 & 1 & \dots & 1 \\ 1 & \omega_1 & \dots & \omega_{n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & \omega_1^{n-1} & \dots & \omega_{n-1}^{n-1} \end{pmatrix}$$

then $\Lambda = \text{diag}(1, \omega_1, \dots, \omega_{n-1})$

$$J = X \Lambda X^{-1}$$

$C = X g(\Lambda) X^{-1}$, so C is diagonalizable.

<可称为 FFT, DFT>

$$|C| = \prod_{k=1}^{n-1} g(\omega_k) \quad (\text{non-zero cyclic means } g(1) \neq 0)$$

• some properties:

① n dim complex matrix B is

diagonalizable $\Leftrightarrow B$ is similar to a cyclic matrix.

② if C is a cyclic matrix, so does C^* . C^{-1} . (use diagonalize)

$\underline{L} \underline{D} \underline{U}$ $\underline{L} \underline{D} \underline{U}$
(triangular / Doolite / Crout decomposition)

2.6 elimination = factorization: $A = LU$

A_{ik} upper left submat.
 $PA = LU$ (or $A = L'P'U$)

the whole forward elimination process (with no row exchanges) is inverted by L , which is still lower triangular, every multiplier l_{ij} stands explicitly at r_{ij} .

and to unitize, pops D (a diag) out.

uniqueness: $L_1 D_1 U_1 = L_2 D_2 U_2$

solving a tri system: $O(\frac{n^2}{2})$
(backward)

(L itself isn't unique) leads to $L_1^{-1} L_2 D_2 = D_1 U_1 U_2^{-1}$

$L_1 = U_1$, must be diag!

elimination / solving $A\vec{x} = \vec{b}$: $O(\frac{n^3}{3})$

so $L_1^{-1} L_2 = I$, $L_1 = L_2 \dots$

o again: echelon! no interact.

o' there also can be written

(\times row 3) row 3 of $U =$ row 3 of $A - l_{31}$ row 1 of $U - l_{32}$ row 2 of U

$A = LU'D$, depends on how you pop D or how

that's row 3 of $A = (l_{31} \ l_{32} \ 1) \cdot \begin{pmatrix} \text{row 1 of } U \\ \text{row 2 of } U \\ \text{row 3 of } U \end{pmatrix}$

U' acts (r/c?)

the algorithm of $A\vec{x} = \vec{b}$: 1. factor into $A = LU$ (multiplier stored in L)

$\frac{n^3}{3}$ n^2

2. solve $Lc = b$ ($b \xrightarrow{L^{-1}} c$)

$Ux = c$ (backward substitution)

n^2 $2nw$

sparse: faster. for a random/full mt. $n=1000$ on a PC takes 1 sec.

$$\Delta \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 2 & 1 & 0 \\ 0 & 1 & 2 & 1 \\ 0 & 0 & 1 & 2 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \end{pmatrix}$$

the zeros at start (r or c) will still be reserved in L or U . then n^3 increases

w is the band width.

□ Cholesky factorization (sq. root factorization) $A = LL^T$, when A is sym. and positive.

(lemma: B is triang., $A = BB^T$ is sym. pos. for $\vec{x}^T A \vec{x} = (\vec{x}B)^T (\vec{x}B) > 0$)

$A = LDU = A^T = U^T D L^T \rightarrow L = U^T$ (B 's row axes)

$A = (L\sqrt{D})(L\sqrt{D})^T$ (he chooses the largest pivot in col in algorithm.)

compare: Lu $n^2 + (n-1)^2 + \dots = \frac{n^3}{3}$

Chol $\frac{n^2}{2} + \frac{(n-1)^2}{2} + \dots = \frac{n^3}{6}$ half \leftarrow sym.'s credit.

algo: $A = \begin{pmatrix} a_{11} & A_1^T \\ A_1 & A' \end{pmatrix} = \begin{pmatrix} l_{11} & \\ & L' \end{pmatrix} \begin{pmatrix} l_{11} & L_1^T \\ & L'^T \end{pmatrix}$, $l_{11} = \sqrt{a_{11}}$, $L_1 = \frac{1}{l_{11}} A_1$

to reduce error since $LUx = b$ not stable (big swallow small)

$L'L'^T = A' - L_1 L_1^T$ — recursion.

$O(\frac{n^2}{2})$ (end: \sqrt{A} , that's $\sqrt{a_{nn}}$)
 L sym.

2.7 transposes and permutations

$$(Ax)^T = x^T A^T, (AB)^T = B^T A^T, (A^{-1})^T = (A^T)^{-1} \text{ note as } A^{-T}$$

inner product. $x \cdot y = x^T y = y^T x$ the idea behind A^T is $Ax \cdot y = x \cdot A^T y$

symmetric matrix $S^T = S$. antisymmetric matrix $S^T = -S$ (skew-sym.)

and $A^T A$ always sym.

orthogonal matrix $Q^T = Q^{-1}$ ($Q Q^T = I$)

(and lu is $S = LDL^T$)

it's standard (unit/orthonormal)

cols of Q are ortho. unitvectors. (orthoaxes)

permu. mt. P puts x_1, \dots, x_n in new order, and $P^T = P^{-1}$, and can be combed by

($n!$ kinds, half odd and half even)

several P_{ij} .

3. Vector Spaces and Subspaces

$$\begin{pmatrix} 0 & -1 \\ 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{matrix} 312 \\ 231 \end{matrix}$$

$$(P_{ij} = P_{ij}^T = P_{ij}^{-1})$$

3.1 spaces of vectors

Vector Space $S \leftarrow$ vec can be matrices even funcs of x .

\mathbb{R}^n contains all real col. vec with n components. if $v, w \in S$, every comb $cv + dw$

Subspace of \mathbb{R}^n is a vecspace inside \mathbb{R}^n , closure to adds and must be in S .

(Z : one-point space consists of $x=0$) must contain zero. nulls)

col space of A contains all combs of A 's cols, which is a subspace of \mathbb{R}^n .

col space contains all vecs Ax . so $Ax=b$ is solvable when b is in $\mathcal{C}(A)$

we call it's spanned by A 's cols.

$\mathcal{C}(A)$: Im A (image) $\mathcal{N}(A)$: Ker A (kernel)

3.2 the nullspace of A : solving $Ax=0$ and $Rx=0$

null space $\mathcal{N}(A)$ (a subspace in \mathbb{R}^n) contains all solution x to $Ax=0$. and it must contain

elimination from A to U to R does not change the nullspace $\mathcal{N}(A) = \mathcal{N}(U) = \mathcal{N}(R)$ Z (zero).

$R = \text{ref}(A)$ means reduced row echelon form. has all pivots = 1, with zeros above and

below. numbers of pivots = numbers of nonzero rows in $R = \text{rank rows (also rank of cols)}$

every matrix with $m < n$ has nonzero

There are $n-r$ free cols. their combs is the

so we have $\dim \mathcal{C}(A) + \dim \mathcal{N}(A) = n$.

solutions in its \mathcal{N} .

(no pivot in) (the comb to produce itself, or $x_j=1$)

(for R (pivot rows and cols) contains I , $x \rightarrow \begin{pmatrix} | & | & | \\ | & | & | \\ | & | & | \end{pmatrix}$)

complete solution of $Ax=0$ or $\mathcal{N}(A)$ ($Rx=0$)

The rank of A is the true size of A or a linear system.

A and U and R have r ind. rows/cols.

full-rank factor: $A = (\text{pivot cols of } A) (\text{first } r \text{ rows of } A) \begin{matrix} m \times n \\ (m \times r) (r \times n) \end{matrix}$

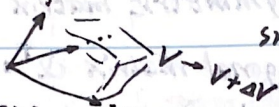
'special solution'
 $S =$
- free colvec
with $x_j=1$ and
zeros.

• 几何理解 LA?

1. basic: 基, 变换, 特征向量, 行列式, Cramer's rule + SVD (works)

2. det and tr: $\det = \prod \lambda$, $\text{tr} = \sum \lambda$, $\text{tr} = \det'$

① $y' = Ay$, $y = e^{tA} y_0$ in small period, trace is the velocity of volume change.
 $|e^{tA}| = 1 + t \cdot \text{tr}(A) + O(t^2)$ small changes mainly in the edges, on its own direction.



② intrinsic (coordinate free):
 mt. is depicted wholly by traces, $\text{tr}(A) + \dots + \text{tr}(A^n) = \lambda_1^n + \dots + \lambda_n^n$.

③ fix-point: $\text{tr} \phi = \sum_i \langle \phi e_i, e_i \rangle$ (like 'how much remained?')
 esp. $A^2 = A$, the rank of $A = A$'s dim of fix-point spanned space = $\text{tr} A$
 for full fix, $\langle \phi, e_i \rangle = 1$.

3. transpose: $A^T x$: x 向列向量基上投影 (实为点乘), 注意与逆区分.

用此解释: $\langle x, Ay \rangle = \langle A^T x, y \rangle$ $\circ (AB)^T = B^T A^T$ $\circ Q^T = Q^{-1}$, $Q^T Q = I$ (not unit: D^2)

• 几个证明与理解:

1. $r_{\text{row}} = r_{\text{col}}$

full-rank decomp.

rref. pivots

2. $|A| = |A^T|$

W^r

suppose $r_{\text{row}} < r_{\text{col}}$.

then the remaining cols ($> r$) could be expressed by W (rank = r)

$A = C \cdot B$
 $m \times n \quad m \times r \quad r \times n$

取 r_{col} 组成 C , 用 B 表达.
 反过来 r 行被 C 组为 A .

$r_{\text{row}}(A) \leq r \leq r_{\text{col}}(A)$.

同理 $r_{\text{col}}(A) \leq r_{\text{row}}(A)$. 故 equal.

geometry analogy



LU/rref 理解

SVD 理解

全排列

定义

symmetric:



since $\langle \bar{x}_1, A \bar{x}_2 \rangle = \langle A \bar{x}_1, \bar{x}_2 \rangle$, if \bar{x}_2 is eigen direction, then $\bar{x}_1 \cdot \bar{x}_2 = 0 \Rightarrow A \bar{x}_1 \cdot A \bar{x}_2 = \lambda^2 \bar{x}_1 \cdot \bar{x}_2 = 0$.

so A maintains the ortho of eigen direction. sym. mt's eigen directions are ortho basis,

(\bar{x}_1 also is) its diagonalize just rotate the axes.

(if $\lambda_1 = \lambda_2$, it's true that the whole plane is eigen.) so $S = Q A Q^T$ orthonormal

• 更多关于 R 的:

R could be write as $\begin{pmatrix} I & F \\ 0 & 0 \end{pmatrix} P$.

'row-col reduced form' $R' = (\text{rref}(R^T))^T = \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix}$

• $A^T A$ always has the same \mathcal{N} with A , ($A^T A x = 0$ 与 $A x = 0$ 为同解方程组)

proof: $A x_0 = 0 \Rightarrow A^T (A x_0) = 0$

$A^T A x_0 = 0 \Rightarrow x_0^T A^T A x_0 = 0, (A x_0)^T (A x_0) = 0, A x_0 = 0.$

② $\% \text{is}$. 正交性或写出 R 矩阵.

($C(A) \perp \mathcal{N}(A^T)$).

so $A^T A x = 0$ only when $A x = 0$

▷ Generably, $r(A^T A) = r(A) (= r(A A^T)) (= r(A) = r(\Sigma))$

3.3 the complete solution to $Ax=b$

o again: If a row/col contains a pivot, it's not a comb of previous rows/cols, otherwise is.
 (with \vec{u}) U tells us which row/col are combs of earlier ones, R tells us what those combs are and the special solution to $Ax=0$.
 R is always the same (determined by A).

Three fundamental subspaces: $\mathcal{C}(A)$ — choose the pivot cols of A as a basis.

Counting Theorem: $\mathcal{R}(A)$ — choose the nonzero rows of R as a basis.

$n = r + n-r$ $\mathcal{N}(A)$ — choose the special solution to $Rx=0$ ($Ax=0$)
 pivot/ind. vars free vars

$(AI) \rightarrow (R, E)$ will virtually tell you everything about A , when A is sq. and inv.,
 R is I and E is A^{-1} .

pivot vars are determined after free vars are chosen.

Complete solution to $Ax=b$: $x = (\text{one particular solution } x_p) + (\text{any } x_n \text{ in nullspace})$
 $(A\vec{b}) \rightarrow (R\vec{d})$ (thus comb of special solution basis)

$Ax=b / Rx=d$ is solvable only when all zero rows of R have zeros in \vec{d} .

full row rank $r=m$ when its col space $\mathcal{C}(A)$ is \mathbb{R}^m , $Ax=b$ is always solvable, but may many.

full col rank $r=n$ when its nullspace $\mathcal{N}(A) = \{0\}$. no free vars, one solution or none.

When is solvable, one particular x_p can be all free vars equal to zero, pivot vars from \vec{d} .

To conclude, there are four cases:

ranks	$r=m=n$	$r=m < n$	$r=n < m$	$r < m, r < n$
types of R	(I)	(I F)	$\begin{pmatrix} I \\ 0 \end{pmatrix}$	$\begin{pmatrix} I F \\ 0 0 \end{pmatrix}$
conclusion	A is inv.	every $Ax=b$ is solvable, but inf.	$Ax=b$ has 1 or 0 solutions	0 or ∞ solutions

o' if $Ax=b$ and $Cx=b$ have same solutions for every b , that means $A=C$.

(proof: let $x = (1, 0, 0, \dots)$ to get col vec $|_A = c$)

3.4 independence, basis and dimension

linear independence: $x_1\vec{v}_1 + x_2\vec{v}_2 + \dots + x_n\vec{v}_n = 0$ only happens when all x 's are zero.

Any set of n vecs in \mathbb{R}^m must be dependent if $n > m$.

basis: (linearly) independent vectors that span the space, like pivot cols of A for $\mathcal{C}(A)$.
 every vector in the space is a unique comb of the basis vecs. (empty set for \mathcal{Z})

dimension: All bases for a space have the same number of vecs called dim, like r_A for $\mathcal{C}(A)$.

Basis of \vec{v}_i 's from basis of \vec{w}_i 's when the change of basis matrix is inv. $V = WB$

dim of outputs + dim of nullspace = dim of inputs. $\Delta \dim(V+W) + \dim(V \cap W) = \dim V + \dim W$

o' $\vec{x} = (x_1, x_2, \dots, x_n)$ 的全排列 $(n!)$ 张成的 S , $\dim S$ 有回钟: $\textcircled{0} \textcircled{0} \textcircled{1} \textcircled{n-1} \textcircled{n}$

($\textcircled{0}$: find a vec that dot to zero / always perpendicular to \vec{x} .)

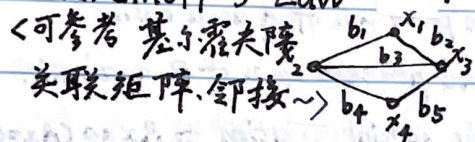
o AB 's rows are combs of B 's rows, AB 's cols are combs of A 's cols.

(so its row/col space contain in B/A 's) $\text{rank}(AB) \leq \text{rank}(A), \text{rank}(B)$. (矩阵乘法的秩)

If multiply by an inv. mt, the rank will not change, because rank can't jump back

so if $AB=0$, we have $C(B) \subset N(A) / C(A^T) \subset N(B^T)$, $r_A + r_B \leq n$. (A: $m \times n$, B: $n \times k$)

o Kirchoff's Law graphs and incidence matrix:



$$Ax=b, A = \begin{pmatrix} -1 & 1 & & & \\ -1 & & 1 & & \\ & -1 & & 1 & \\ & & -1 & & 1 \\ & & & -1 & 1 \end{pmatrix}$$

since $\vec{b}=0 \Rightarrow x_1=x_2=x_3=x_4$
thus $\vec{x} = (c, c, c, c)$ in $N(A)$,
the rank must be 3.

since $b_1, -b_2, -b_3$ is a loop, rows 1, 2, 3 of A
are not lin. ind. but $(1, -1, 1, 0, 0)$ will be in A 's left nullspace.

(two ind. loops, $m-r=5-3=2$)

o Kirchoff's Voltage Law $Ax=b$

Kirchoff's Current Law $A^T y=0$ \rightarrow 'balance equation': most important eq. in applied maths.
(y_i flows on b_i edge)

o $A=UV^T$'s four subspaces are $\begin{matrix} \swarrow & \searrow \\ v & v^\perp \\ \nwarrow & \nearrow \\ u & u^\perp \end{matrix}$. If B have those same subspaces, then $B=CA$.
Generally, mt. A and B have same four subspaces means their rref equals.

(proof: same row space is the key. every row in $\text{rref}(A)$ must be comb of $\text{rref}(B)$, but notice I in these rref, it means each row should be equal in rref.)

o More Advanced Level: Matrix Spaces and Function Spaces

o dim of the subspace of sym. mts. is $\frac{1}{2}n^2 + \frac{1}{2}n$.

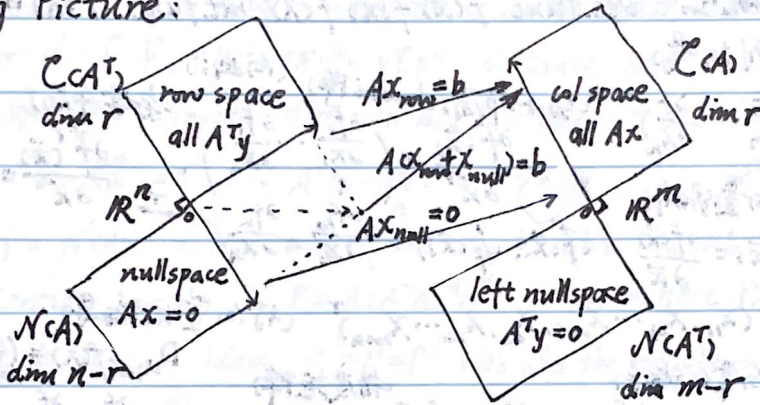
o second derivative equations: solution space $y'' = -y$ has two basis funcs: $\sin x$ and $\cos x$.
but $y'' = 2$ don't form a subspace,
a particular solution $y_p = x^2$. then the complete solution is $y = x^2 + cx + d$.
 $y'' = y$: e^x and e^{-x} .
 $y'' = 0$: x and 1 .

Ex. $A = \begin{pmatrix} 1 & 0 & -1 \\ -1 & 1 & 0 \\ 0 & -1 & 1 \end{pmatrix}$, notice $A \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$. The counting theorem writes $6 + 3 = 9$.
 $AX=B, X \in 3 \times 3$ mt. space.
 $C \ N$

o Fredholm's Alternative: Exactly one of these problems has a solution $\begin{cases} Ax=b \\ A^T y=0 \text{ with } y^T b \neq 0 \end{cases}$

3.5 dimensions of the four subspaces

The Big Picture:



1. A has the same row space as R . same dim r and same basis.
2. The col space of A has dim r . col rank equals row rank. (Ranking Theorem) $\mathcal{C}(A) \neq \mathcal{C}(R)$! but dim r . (for same combs of col = 0 for A and R , thus $Ax=0 \Leftrightarrow Rx=0$) ($\mathcal{N}(A) = \mathcal{N}(R)$)
3. A has same nullspace as R . same dim $n-r$ and same basis. (again, doesn't change solutions)
4. The left nullspace of A has dim $m-r$. (outgoing Theorem or $R^T y=0 / y^T R=0$, $R=EA$)

Fundamental Theorem of Linear Algebra:

Part 1. The col space and row space both have dim r ; The nullspaces have dim $n-r$ and $m-r$.

rank-one matrix: $A = uv^T = \text{col} \times \text{row}$. $\mathcal{C}(A)$ has basis u , $\mathcal{C}(A^T)$ has basis v .

rank-one's sum: $EA=R$, $A=E^{-1}R=CR$ Ex. $A = \begin{pmatrix} 1 & 0 & 3 \\ 1 & 1 & 7 \\ 4 & 2 & 20 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 4 & 2 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 3 \\ 0 & 1 & 4 \\ 0 & 0 & 0 \end{pmatrix} = CR$.

(pre-read: a full-rank decomp.) Every rank r mat. is a sum of r rank-one mts: \circ left eigenvcs.

(rank-one expansion)

$$A = (\vec{u}_1 \ \vec{u}_2 \ \vec{u}_3) \begin{pmatrix} \vec{v}_1^T \\ \vec{v}_2^T \\ \text{zero rows} \end{pmatrix} = \vec{u}_1 \vec{v}_1^T + \vec{u}_2 \vec{v}_2^T = \text{rank 1} + \text{rank 1}$$

4. Orthogonality

4.1 orthogonality of the four subspaces

orthogonal vectors $v^T w = 0$ and have $\|v\|^2 + \|w\|^2 = \|v+w\|^2 = \|v-w\|^2$
 The orthogonal complement of a subspace V contains every vector that is perpendicular to V . denoted by V^\perp
 Orthogonality is impossible when all v in V and all w in W .

Part 2. $\mathcal{N}(A)$ is the ortho complement of $\mathcal{C}(A^T)$; (in \mathbb{R}^n) $\dim V + \dim W > \dim(\text{whole space})$
 $\mathcal{N}(A^T)$ is the ortho complement of $\mathcal{C}(A)$. (in \mathbb{R}^m) (since it only can be zero to be contained in both ortho spaces.)

This means every x can be split into $x_r + x_n$. When $v^T v = 0 \Rightarrow v = 0$, a row of A can't be in $\mathcal{N}(A)$.
 A multiplies, nullspace component x_n goes to zero $Ax_n = 0$.
 row space component x_r goes to colspace $Ax_r = Ax$.

Every b in colspace comes from one and only one vec x_r in row space.

There is an $r \times r$ inv. mt. hiding inside A , if throw away two nullspaces.

From the row space to the colspace, A is inv. ('pseudoinverse', see 7.4) n vecs

\circ the definition of basis has two properties, but one implies the other: ind. $\Leftrightarrow \text{span } \mathbb{R}^n$

更多关于 tri. mt. 的: L/U 的 inv 仍为 L/U; $L_1, L_2 / U_1, U_2$ 积仍为 L/U; $LX = L' / UX = U'$ 列 (且对角线 $d = d_1 d_2$) X 一定也为 L/U

matrix derivative

(向量值函数)

矩阵求导

定义概念 vec. func. $\vec{f}(x)$ $\vec{f}(x)$ $\vec{f}(x)$ mt. func. $F(x) = (F_{ij}(x_{mn}))$

(标函):

向量: $\frac{df(x)}{dx} = \begin{pmatrix} \frac{df}{dx_1} \\ \vdots \\ \frac{df}{dx_n} \end{pmatrix}$ (果求列向量) $\frac{df(x)}{dx^T} = \left(\frac{df}{dx_1} \dots \frac{df}{dx_n} \right)$ (分子布局, Jacobi 阵) $\frac{df(x)}{dx^T} = \begin{pmatrix} \frac{df_1}{dx_1} \dots \frac{df_1}{dx_n} \\ \vdots \\ \frac{df_m}{dx_1} \dots \frac{df_m}{dx_n} \end{pmatrix}$ (向函: 分母布局) $\frac{df^T(x)}{dx} = \begin{pmatrix} \frac{df_1}{dx_1} \dots \frac{df_m}{dx_1} \\ \vdots \\ \frac{df_1}{dx_n} \dots \frac{df_m}{dx_n} \end{pmatrix}$

(行偏导)

向量 $D_x f(x) = \frac{df(x)}{dx^T} \xrightarrow{I} \nabla_x f(x) = \frac{df(x)}{dx}$ ($\vec{f}(x), \vec{f}(x)$) $\frac{df_1}{dx_1} \dots \frac{df_m}{dx_n}$

矩阵: $\text{vec}(X)$: 列堆栈 向量化 $= (x_{11}, x_{21}, \dots, x_{m1}, x_{12}, x_{22}, \dots, x_{mn})^T$ (行向量化偏导)

(Jacobi 阵) $D_x f(x) = \begin{pmatrix} \frac{df}{dx_{11}} \dots \frac{df}{dx_{m1}} \\ \vdots \\ \frac{df}{dx_{1n}} \dots \frac{df}{dx_{mn}} \end{pmatrix}$ (列向量化偏导) $\nabla_{\text{vec}} f(x) = \frac{df(x)}{d\text{vec}} = \left(\frac{df}{dx_{11}} \dots \frac{df}{dx_{mn}} \right)^T$ (梯度矩阵) $\nabla_x f(x) = \begin{pmatrix} \frac{df}{dx_{11}} \dots \frac{df}{dx_{m1}} \\ \vdots \\ \frac{df}{dx_{1n}} \dots \frac{df}{dx_{mn}} \end{pmatrix}$ $D_{\text{vec}} f(x) = \left(\frac{df}{dx_{11}}, \frac{df}{dx_{21}}, \dots, \frac{df}{dx_{mn}} \right)$

同理矩阵也有

$D_x F(x) = \frac{\partial \text{vec}(F(x))}{\partial \text{vec} X}$ $\nabla_x F(x) = \frac{\partial \text{vec}^T(F(x))}{\partial \text{vec} X}$ Δ Transpose 的 T 位置可以变动.

向量标函: $\frac{\partial c}{\partial x} = 0$ $\frac{\partial f(x)g(x)}{\partial x} = \frac{df(x)}{dx} g(x) + f(x) \frac{dg(x)}{dx}$ 等其他老规律, 常用新规律:

1. $\frac{\partial (x^T a)}{\partial x} = \frac{\partial (a^T x)}{\partial x} = a$ 2. $\frac{\partial (x^T x)}{\partial x} = 2x$ 3. $\frac{\partial (x^T A x)}{\partial x} = Ax + A^T x$ 4. $\frac{\partial (a^T x x^T b)}{\partial x} = (ab^T + ba^T)x$

矩阵标函: 5. $\frac{\partial (a^T X b)}{\partial X} = ab^T$ 6. $\frac{\partial (a^T X^T b)}{\partial X} = ba^T$ 7. $\frac{\partial (a^T X X^T b)}{\partial X} = (ab^T + ba^T)X$ 8. $\frac{\partial (a^T X^T X b)}{\partial X} = X(ab^T + ba^T)$

矩阵标函 $df(x) = \frac{df}{dx_{11}} dx_{11} + \dots + \frac{df}{dx_{mn}} dx_{mn} \triangleq \text{tr} \left(\frac{df(x)}{dx^T} dx \right)$ 常用新规律:

1. $d(AXB) = A dX B$ 2. $d|X| = |X| \text{tr}(X^{-1} dX) = \text{tr}(X^* dX)$ 3. $dX^{-1} = -X^{-1} dX X^{-1}$ ($XX^{-1} = I$ 作微分)

$\Delta df(x) = d(\text{tr} f(x)) = \text{tr}(df(x))$ (标量不变)

例: 1. $\frac{\partial (a^T X X^T b)}{\partial X}$: $d(a^T X X^T b) = \text{tr}(a^T d(X X^T) b) = \text{tr}(a^T dX X^T b) + \text{tr}(a^T X dX^T b)$ ($d(X^T) = (dX)^T$)

故 $\frac{\partial (a^T X X^T b)}{\partial X^T} = X(ba^T + ab^T)$, $\text{tr}(AB) = \text{tr}(BA)$ $\text{tr}(X^T b a^T dX) + \text{tr}(X^T a b^T dX)$

2. $\frac{\partial \text{tr}(X^T X)}{\partial X} = 2X$: $d(\text{tr}(X^T X)) = \text{tr}(dX^T X) + \text{tr}(X^T dX) = 2\text{tr}(X^T dX)$

3. $\frac{\partial \log|X|}{\partial X} = X^{-T}$: $d(\ln|X|) = \text{tr} \left(\frac{1}{|X|} d|X| \right) = \text{tr}(X^{-1} dX)$ 故 $\frac{\partial \text{tr}(X^T X)}{\partial X^T} = 2X^T$, $\frac{\partial \text{tr}(X^T X)}{\partial X} = 2X$

5. $\frac{\partial |X|^3}{\partial X} = 3|X|^2 X^{-T}$ 6. 概统估计: $x \sim N(\mu, \Sigma)$ 4. $\frac{\partial \text{tr}(X+A)^{-1}}{\partial X} = -(X+A)^{-2} X^T$

极大似然 $\ln L(\mu, \Sigma) = \sum \ln f(x_i) = -\frac{p}{2} n \ln 2\pi - \frac{1}{2} \sum (x_i - \mu)^T \Sigma^{-1} (x_i - \mu)$
 有 $\frac{\partial \ln L(\mu, \Sigma)}{\partial \mu} = \sum^{-1} \sum \frac{\partial}{\partial \mu} (x_i - \mu) = 0$
 $\left\{ \begin{aligned} \frac{\partial \ln L(\mu, \Sigma)}{\partial \mu} &= 0 \\ \frac{\partial \ln L(\mu, \Sigma)}{\partial \Sigma} &= 0 \end{aligned} \right.$ (充分)

4.2 projections

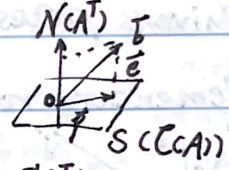
(m dim.)

The projection of \vec{b} onto a subspace S is the closest vector \vec{p} in S with $\vec{b}-\vec{p}$ is ortho to S .

error: $\vec{e} = \vec{b} - \vec{p}$. $\|\vec{p}\|^2 + \|\vec{e}\|^2 = \|\vec{b}\|^2$. distance: $\|\vec{e}\|$

A 's cols span S , suppose they are ind. (otherwise use those r ind. cols)

$\vec{a}_i \cdot (\vec{b} - A\vec{x}) = 0 \rightarrow A^T(\vec{b} - A\vec{x}) = 0, p = A\vec{x}$. $A^T\vec{b} = A^T A\vec{x}$



$N(CA^T) = N(CA)$: $A^T A$ is inv. if and only if A has ind. cols. so $p = A(CA^T A)^{-1} A^T b$.

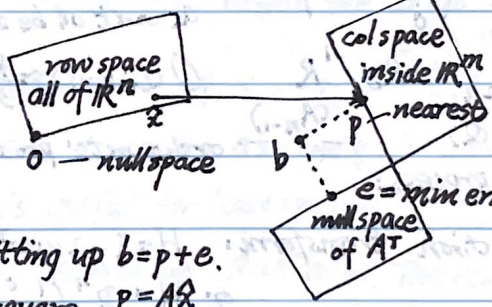
The projection matrix is $P = A(CA^T A)^{-1} A^T$ Δ we have $P_C A P_R = A (= P_C A = A P_R)$. $P_{C/R}$: proj. to col/row space.

proj. mt is sym. sq. idem. $P^2 = P = P^T$ (it's also the definition of proj. mt.) $\text{rank}(P) = \text{rank}(A) = r$

(for a vec, $\hat{x} = \frac{a^T b}{a^T a}$, $p = \frac{a a^T}{a^T a}$ a rank-one mt.) (When A is inv., P is I) ($P \sim (I_r \ 0)$)

4.3 least squares approximations

$A = \begin{pmatrix} t_1 \\ \vdots \\ t_m \end{pmatrix}$ $x = \begin{pmatrix} C \\ D \end{pmatrix}$ $b = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix}$
 $(b = C + Dt, m > 2, n = 2)$



Δ since first col of A is $(1, 1, \dots, 1)$, $\sum e_i = 0$ or t mean will go to b mean. ($\vec{b} = C + D\vec{e}$)

There are no solutions to $Ax = b$, we splitting up $b = p + e$.

\hat{x} is the least square solution, for the square $p = A\hat{x}$

length of $Ax - b$ is minimized: $E = \|Ax - b\|^2$ min. ($= \sum e_i^2$)

By geometry: the nearest point is the projection. solve $A^T A x = A^T b$.

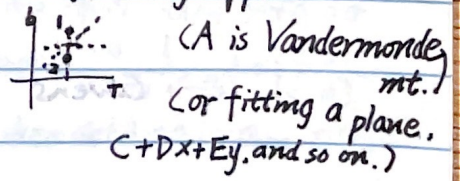
By algebra/calculus: the partial derivatives of E are zero. $\frac{\partial E}{\partial C} = 0, \frac{\partial E}{\partial D} = 0$ or $\frac{\partial E}{\partial x} = 0$ as a whole, we get $(Ax - b)^T A = 0$.

$A^T A = \begin{pmatrix} m & \sum t_i \\ \sum t_i & \sum t_i^2 \end{pmatrix}$ $A^T b = \begin{pmatrix} \sum b_i \\ \sum t_i b_i \end{pmatrix}$, $\hat{x} = \begin{pmatrix} C \\ D \end{pmatrix} = (A^T A)^{-1} A^T b$

If $\sum t_i = 0$ then A has ortho cols and $A^T A$ is diag. It's worth shifting times t_i by subtracting the average $\hat{t} = \frac{\sum t_i}{m}$, $T = t - \hat{t}$ or make cols ortho. in advance by Gram-Schmidt. $b = C + D(t - \hat{t})$

Like fitting a straight line, the unknowns $C + Dt + Et^2$ still appear linearly in fitting a parabola.

(if A has dependent cols, we have many solution lines, Then $\hat{x} = (C, D, E)$, nothing diff. pseudo inv. of A will choose the shortest solution to $A\hat{x} = p$. when A is ind. cols the pseudo inv. is usual left inv. $L = (A^T A)^{-1} A^T$)



(A is Vandermonde mt.) (or fitting a plane, $C + Dx + Ey$, and so on.)

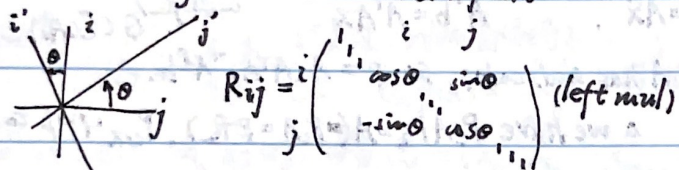
• $A = QR = QHR'$ (diag R all 1s. with heights h_i in h_i tells the height col i above the plane of H)
 $O(n) \approx n^3$

□ ways of QR decomp.: Householder decomp. and Givens decomp.

QR分解之豪斯霍尔德分解, 吉文斯分解

• Givens process: (A has ind. cols!)

elementary rotation transform:



• QR's uniqueness: $Q_1 R_1 = Q_2 R_2$

$Q_1 = Q_2 R_2 R_1^{-1} = Q_2 U$ leads to

$Q_1^T Q_1 = I = U^T U$ must be diag!

since we make R positive diag, so $U = I$. $R_1 = R_2 \dots$

choose $\cos\theta = \frac{x_i}{\sqrt{x_i^2 + x_j^2}}$ rotate x in $i-j$ plane to $x_j' = 0$. $x_i' = \sqrt{x_i^2 + x_j^2}$.
 notice that R_{ij} is orthogonal.

$A_1 = R_{1n} R_{1n-1} \dots R_{12} A = \begin{pmatrix} \|a_1\| & & & \\ 0 & \dots & & \\ 0 & & \dots & \\ \vdots & & & \end{pmatrix}$

then rotate a_2 .
 (since A ind. cols, a_2 can't all be zeros.)

$A_2 = R_{2n} R_{2n-1} \dots R_{23} A_1 = \begin{pmatrix} \|a_1\| & a_{21} & & \\ 0 & \|a_2 - a_{21}\| & \dots & \\ \vdots & 0 & \dots & \\ 0 & 0 & & \end{pmatrix}$

so $A = (R_{n-1} R_{n-2} \dots R_{12})^{-1} R$
 $(Q) \quad (A_{n-1})$

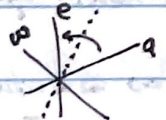
$O(\text{Givens}) \approx \frac{4}{3}n^3$ it's better for sparse matrix.

• Householder process:

elementary reflection transform:

$H = I - 2\omega\omega^T$ ω is mirror normal unit vec.
 notice that H is orthogonal and sym.
 or $H = Q^{-1} \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} Q$ ($|H| = -1$, $|R_{ij}| = 1$)
 where $Q\omega = \begin{pmatrix} 0 \\ \vdots \\ 1 \end{pmatrix}$ ($H = 2P - I$ ano. way)

choose $\omega = \frac{a - \|a\|e}{\|a - \|a\|e\|}$ to reflect a to e



$H_1 A = \begin{pmatrix} \|a_1\| & \dots \\ 0 & \dots \\ \vdots & \vdots \\ 0 & \dots \end{pmatrix}$

then preserve col 1 and row 1.

find $H_2 A_2$ to make $\begin{pmatrix} a_{22} & \dots \\ 0 & \dots \\ \vdots & \vdots \\ 0 & \dots \end{pmatrix}$. $H_2 = \begin{pmatrix} 1 & 0 \\ 0 & H_2' \end{pmatrix}$

choose $e_i = (1, 0, \dots, 0)^T$

so $A = (H_{n-1} \dots H_1)^{-1} R$
 $(Q) \quad (A_{n-1})$ $O(\text{House}) \approx \frac{2}{3}n^3$

To conclude:

- Gram-Smidt small mt.; generate q in every step
- not for sparse mt.; not stable; save all mt. thus cost large space.
- Householder suitable for dense mt.; not calculate Q explicitly (if only need R)
- not generate q 's until finished; not for sparse mt.
- Givens suitable for sparse mt.
- not generate q 's until finished; not for dense mt.

• a rotation mt. = two reflection mts: $R_{ij} = H_1 H_2$

$H_1 = I - 2\omega_1\omega_1^T$, $\omega_1 = (0, \dots, \sin\frac{\theta}{2}, 0, \dots, \cos\frac{\theta}{2}, 0, \dots, 0)^T$; $H_2 = I - 2\omega_2\omega_2^T$, $\omega_2 = (0, \dots, \sin\frac{\theta}{2}, 0, \dots, \cos\frac{\theta}{2}, 0, \dots, 0)^T$

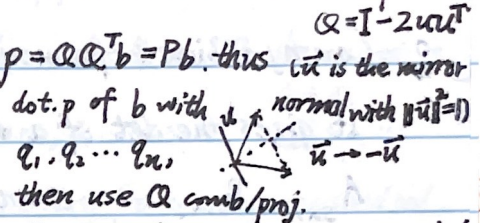
4.4 orthonormal bases and Gram-Schmidt

not require sq!

A matrix Q with orthonormal cols satisfies $Q^T Q = I$. It leaves lengths unchanged: $\|Qx\| = \|x\|$ and also preserves dot.p $(Qx)^T (Qy) = x^T y$. Some Q 's examples: Rotation, Permutation, Reflection

The least squares solution to $Qx = b$ is $\hat{x} = Q^T b$, the proj: $p = QQ^T b = Pb$. thus \hat{x} is the mirror (If Q is sq./orthogonal, it spans all space so $P = I$. $Q^T = Q^{-1}$) dot.p of b with q_i normal with $\|q_i\|=1$

Gram-Schmidt process:



Subtract from every new vec its proj in the

The one-dim proj are uncoupled.

$$\begin{aligned} (A=a) \quad C &= c - \frac{A^T c}{A^T A} A - \frac{B^T c}{B^T B} B / \|C\| \\ B &= b - \frac{A^T b}{A^T A} A / \|B\| \end{aligned}$$

directions already set. $b = \sum q_i (q_i^T b)$

At the end, divide the ortho. vecs by their length.

factorization: $A = QR$ (orthotri. decomp.)

$$(q_1, q_2, \dots, q_n) \begin{pmatrix} q_1^T a_1 & q_1^T a_2 & \dots & q_1^T a_n \\ q_2^T a_1 & q_2^T a_2 & & q_2^T a_n \\ & & & q_n^T a_n \end{pmatrix}$$

$a_1, \dots, a_n \rightarrow q_1, \dots, q_n$ later q 's are ortho. to earlier a 's (and q 's.) $R = Q^T A$ to recognize the upper tri. Any $m \times n$ ind. cols A can be factored. It's useful for least squares: (R 's diag is $\|a_1\|, \|a_2\|, \dots$)

$A^T A = R^T R$. we solve $R\hat{x} = Q^T b$ by back substitution ($\hat{x} = R^{-1} Q^T b$). The real cost is $G-S: O(mn^2)$.

(modified G-S: more stable than original which subtracts all proj at once. it subtracts one proj at a time as in

$$C^* = a_3 - (q_1^T a_3) q_1, \quad C = C^* - (q_2^T C^*) q_2, \quad q_3 = \frac{C}{\|C\|}$$

5. Determinants

ortho. $|Q| = 1$

5.1 the properties of determinants

$|\det A| =$ volume of box whose sides are A 's rows/cols. (sign is the direction.)

⊕ the det of $I_{n \times n}$ is 1. ⊕ the det changes sign when two row exchanged.

⊕ the det is a linear func. of each row separately. $\begin{vmatrix} a+ta' & b+tb' \\ c & d \end{vmatrix} = \begin{vmatrix} a & b \\ c & d \end{vmatrix} + t \begin{vmatrix} a' & b' \\ c & d \end{vmatrix}$

then there are some important corollaries: ⊕ if two rows of A are equal, then $\det A = 0$.

⊕ subtracting a multiple of one row from ano. row leaves A unchanged. $\begin{vmatrix} a & b \\ c-ta & d-tb \end{vmatrix} = \begin{vmatrix} a & b \\ c & d \end{vmatrix}$

⊕ a mt. with a row of zeros has $\det = 0$. ⊕ if A is triangular then $\det A = a_{11} \dots a_{nn}$.

⊕ singular A 's $\det = 0$. inv. A 's $\det \neq 0$. ⊕ the det AB is $\det A$ times $\det B$. esp. $\det A^{-1} = \frac{1}{\det A}$.

⊕ the transpose A^T has the same det as A .

(proof: $\frac{|AB|}{|B|}$ has the same three properties that define $|A|$ (⊕⊗⊕), and the case $|B|=0$ is easy for AB is singular.)

(use $PA = LU$, $A^T P^T = U^T L^T$ and ⊕ proof)

Δ try to image those in geometry!

o' if A is singular, $AC^T = 0$, then each col of C^T is in the nullspace of A .
 ($A_{n \times n}$ of rank $n-2$ or smaller, all cofactors are zero and we only find $\vec{x} = 0$.)

o $AC^T = (\det A)I \rightarrow (\det A)^{n-1} = \det C$. so if you know all cofactors of A (with positions) you could get A .

o Cauchy-Binet formula: 柯西-比内公式

It gives the det of a sq. mt. AB (esp. $A^T A$) when the factors A, B are rectangular.

$A_{s \times n}, B_{n \times s}$: $s > n$ leads to $\det AB = 0$;
 $s \leq n$ then $\det AB = \text{sum of } (s \times s \text{ det in } A)(s \times s \text{ det in } B)$.

$$\begin{pmatrix} a & b & c \\ d & e & f \end{pmatrix} \begin{pmatrix} g & j \\ h & k \\ i & l \end{pmatrix} = \begin{pmatrix} a & b & c \\ d & e & f \end{pmatrix} \begin{pmatrix} g & j \\ h & k \\ i & l \end{pmatrix} = \begin{pmatrix} a & b & c \\ d & e & f \end{pmatrix} \begin{pmatrix} g & j \\ h & k \\ i & l \end{pmatrix} \quad \begin{matrix} (s's \text{ cols of } A \text{ correspond with } s's \text{ rows of } B) \\ \text{(calculate the } 2 \times 2 \text{ det's underlined)} \end{matrix}$$

o Jacobi formula: The cofactor formula can be generalized. (or 'Laplace expansion')

For $n \times n$ mt., we can choose $k \times k$ det times $(n-k) \times (n-k)$ det. Then add them with signs. The sign depends on the permutation of $a_1, \dots, a_k, b_1, \dots, b_{n-k}$ (a and b is the column you choose)

$$\begin{vmatrix} 2 & -1 & 0 & 0 & 0 \\ -1 & 2 & -1 & 0 & 0 \\ 0 & -1 & 2 & -1 & 0 \\ 0 & 0 & -1 & 2 & -1 \\ 0 & 0 & 0 & -1 & 2 \end{vmatrix} = \begin{vmatrix} 2 & -1 \\ -1 & 2 \end{vmatrix} \cdot \begin{vmatrix} 2 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & -1 & 2 \end{vmatrix} \cdot P(1,2,3,4,5) + \begin{vmatrix} 2 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & -1 & 2 \end{vmatrix} \cdot \begin{vmatrix} -1 & -1 & 0 \\ 0 & 2 & -1 \\ 0 & -1 & 2 \end{vmatrix} \cdot P(1,3,2,4,5) = 6$$

o rank-one mt. $A = uv^T$ has an eigenvector u . its $\lambda = v^T u$. Other eigenvalues are zeros. (eigenvecs on the plane v^\perp)

更多例子
 特征值: permutation mt. P have $P^k = I \Rightarrow \lambda^k = 1$. Then we find P 's eigens. (maybe complex)

$\sum \lambda_i^k = \text{tr } A^k$. (every perm. mt. leaves $\vec{x} = (1, 1, \dots, 1)^T$ unchanged so $\lambda_1 = 1$)

rotation mt. can do the same. $\lambda^{\frac{2\pi}{\theta}} = 1$ we find $e^{\pm i\theta}$ (for higher dim, $\lambda = 1, 1, \dots, e^{\pm i\theta_{ij}}$)

o' When a mt. is diag, the number of nonzero eigenvalues is its rank. $\sim \begin{pmatrix} \Lambda_r \\ 0 \end{pmatrix}$

o $A+B$. AB might not have $\lambda = \lambda_a + \lambda_b, \lambda_a \lambda_b$! It only happens when A, B have same eigenvecs.

o esp. $B=I$, $A+cI$ can have $\lambda = \lambda_a + c$ for A 's eigenvecs must in I 's eigen.

(when λ has multiplicity, the whole plane spanned by its eigenvecs are eigen comb; for I it's whole space)

o no interconnection between invertibility diagonalizability!

$$\lambda = 0? \quad X \exists ?$$

5.2 permutations and cofactors

The determinant of any $n \times n$ matrix can be found in three ways: (more: Big Partition...)

1. pivot formula: $PA = LU \rightarrow \det A = \pm \det U = \pm (\text{product of pivots})$

(we don't need row exchanges when all the upper left submts. have $\det A_k \neq 0$. The k th pivot is

2. big formula: $\det A = \text{sum over all } n! \text{ col permutations } P = (\alpha, \beta, \dots, \omega) \quad a_k = \frac{\det A_k}{\det A_{k-1}}$

$= \sum \det P \cdot a_{1\alpha} a_{2\beta} \dots a_{n\omega}$ choose one entry from every row and col.

$$\begin{vmatrix} a_{11} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{31} & \dots & a_{3n} \end{vmatrix} = a_{11} a_{22} a_{33} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix} + a_{12} a_{23} a_{31} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix} + a_{13} a_{21} a_{32} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix} + a_{11} a_{23} a_{32} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix} + a_{12} a_{21} a_{33} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix} + a_{13} a_{22} a_{31} \begin{vmatrix} | & | & | \\ \hline & & \\ \hline \end{vmatrix}$$

$n!$ ways to order, half odd and half even.

3. cofactor formula: $\det A = a_{i1} C_{i1} + \dots + a_{in} C_{in}$. each cofactor C_{ij} includes its correct sign:

since a_{ij} removes row i and col j , we can linearly expand det by r_i/c_j . $C_{ij} = (-1)^{i+j} M_{ij}$

$$\begin{vmatrix} a_{11} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{31} & \dots & a_{3n} \end{vmatrix} = \begin{vmatrix} a_{11} & & \\ a_{22} & a_{23} & \\ a_{32} & a_{33} & \end{vmatrix} + \begin{vmatrix} a_{21} & a_{23} & \\ a_{31} & a_{33} & \end{vmatrix} + \begin{vmatrix} a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & \end{vmatrix}$$

to permute a_{ij} to a_{11} need $(-1)^{i+j-2}$ reverse, then M_{ij} 's det permutes as itself.

$$= a_{11}(a_{22}a_{33} - a_{32}a_{23}) + a_{12}(a_{23}a_{31} - a_{21}a_{33}) + a_{13}(a_{21}a_{32} - a_{22}a_{31}) \quad (\det A^T = \det A, \text{ so did } \text{col } j.)$$

5.3 Cramer's rule, inverse, and volumes

Cramer's rule solves $Ax = b$. we construct $A \begin{pmatrix} x_1 & 0 & 0 \\ x_2 & 1 & 0 \\ x_3 & 0 & 1 \end{pmatrix} = (b \ a_2 \ a_3) = B_1$. $x_1 = \frac{\det B_1}{\det A}$

then $x_n = \frac{\det B_n}{\det A}$, where $\text{mt. } B_n$ has the n th col of replaced by b .

Then we find the cols of A^{-1} by solving $AA^{-1} = I$. $B_n = \begin{pmatrix} \dots & 0 & \dots \\ \vdots & \vdots & \vdots \\ \dots & 1 & \dots \end{pmatrix}$ so

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \quad (A^{-1})_{ij} = \frac{C_{ji}}{\det A} \quad \text{and } A^{-1} = \frac{C^T}{\det A}$$

A direct proof of $AC^T = (\det A)I$:

(C_{ij} go into the cofactor matrix C)

The cofactor formula yields $\det A$ on the diagonal, multiplying cofactors from diff rows

is the cofactor formula for a new matrix when the second row is copied into first row,

so yields zero.

$$\text{The triangle with corners } (x_1, y_1), (x_2, y_2), (x_3, y_3) \text{ has area} = \frac{\det}{2} = \frac{1}{2} \begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}$$

We can prove that the parallelogram's area when $(x_2, y_2) = (0, 0)$, $\text{area} = \frac{1}{2} \begin{vmatrix} x_1 & y_1 \\ x_2 & y_2 \end{vmatrix}$.

or the box's volume obeys the determinant properties ①②③.

cross p.: $\vec{u} \times \vec{v} = \begin{vmatrix} \hat{i} & \hat{j} & \hat{k} \\ u_1 & u_2 & u_3 \\ v_1 & v_2 & v_3 \end{vmatrix}$ is perpendicular to \vec{u} and \vec{v} .

triple p.: $(\vec{u} \times \vec{v}) \cdot \vec{w} = \begin{vmatrix} u_1 & u_2 & u_3 \\ v_1 & v_2 & v_3 \\ w_1 & w_2 & w_3 \end{vmatrix} = \text{determinant} = \text{volume}$.

(antisym.)

• Heisenberg's uncertainty principle: (inf. mt.) position mt. P , momentum mt. Q . (sym.)

$$QP - PQ = I \quad (\text{To have } P\pi = 0 \text{ and } Qx = 0 \text{ would require } x = 0)$$

$$x^T x = x^T QP x - x^T P Q x \leq 2 \|Px\| \|Qx\| \quad \text{Explain that last step by using Schwarz ineq.}$$

$$\|u^T v\| \leq \|u\| \|v\|$$

So $\frac{\|Px\| \|Qx\|}{\|x\|^2} \geq \frac{1}{2}$. (Impossible to get position and momentum error all very small.)

(proof via)

1. A and B share the same n ind. eigenvcs. if and only if $AB = BA$. (commutable)

2. If A is $m \times n$ and B is $n \times m$, then AB and BA have some nonzero eigenvalues.

(and $|n-m|$ zeros.)

$$\text{(proof: } \begin{pmatrix} I & -A \\ & I \end{pmatrix} \begin{pmatrix} AB & 0 \\ B & 0 \end{pmatrix} \begin{pmatrix} I & A \\ & I \end{pmatrix} = \begin{pmatrix} 0 & 0 \\ B & BA \end{pmatrix} \begin{pmatrix} m \times m & m \times n \\ n \times m & n \times n \end{pmatrix})$$

E has m eivalues of $AB + n$ zeros

so $E = \begin{pmatrix} AB & 0 \\ B & 0 \end{pmatrix} \sim F = \begin{pmatrix} 0 & 0 \\ B & BA \end{pmatrix}$, they have some $m+n$ values: F has n eivalues of $BA + m$ zeros

3. Suppose A_1, A_2 $n \times n$ inv., then $A_1 A_2$ is similar to $A_2 A_1$, same eigenvalues. (Choose A_2)

• Gershgorin circles: Every λ is in the circle around one or more diagonal entries a_{ii} :

$$\text{(reason: } A - \lambda I \text{ is singular, it cannot be } |a_{ii} - \lambda| \leq R_i = \sum_{j \neq i} |a_{ij}| \text{ diag dominant.)}$$

• eigenvalues of A equals eigenvalues of A^T . (since $\det(A - \lambda I) = \det(A^T - \lambda I)$)

The eigenvectors of A and A^T from different λ 's are perpendicular, (esp. sym. mt.'s eigenvcs are ortho.)
since $A\vec{x}_1 \cdot \vec{x}_2 = \vec{x}_1 \cdot A^T \vec{x}_2 \Rightarrow \lambda_1 \vec{x}_1 \cdot \vec{x}_2 = \lambda_2 \vec{x}_1 \cdot \vec{x}_2$, $\lambda_1 \neq \lambda_2$ leads to $\vec{x}_1 \perp \vec{x}_2$. (orthog.)

(or $A = X\Lambda X^{-1}$, $A^T = X^{-T}\Lambda X^T$, the eigenvcs of A^T are X^{-T} , $X^{-T}X = I$ leads to $\vec{x}_i \cdot \vec{x}_j$ ($i \neq j$) = 0.) ($n = r(A)$)

• 矩阵收集包 (三): \triangleright left eigenvectors of A ($y^T A = \lambda y^T$). rank-one's sum: $A = \lambda_1 x_1 y_1^T + \dots + \lambda_n x_n y_n^T$.

• Hadamard matrix $H_n = \begin{pmatrix} 1 & 1 \\ & 1 & -1 \end{pmatrix}$ $H_{ij} = \pm 1$ and orthogonal. (so n cannot be odd) $H^T H = nI$, every cols or rows in H are orthogonal. ($H = H^T = H^{-1}$)

(Sylvester theorem) $H \checkmark \rightarrow \begin{pmatrix} H & H \\ H & -H \end{pmatrix} \checkmark$ Hadamard Theorem:

(Hadamard conjuncture: $n | 4$, $H_n \checkmark$)

$$A = (a_{ij})_{n \times n} \text{ with } -1 \leq a_{ij} \leq 1,$$

then $|\det A| \leq n^{\frac{n}{2}}$. equality achieves when A is H ($a_{ij} = \pm 1$ and rows are ortho.)

• Hessenberg matrix $H_3 = \begin{pmatrix} 2 & 1 \\ 1 & 2 & 1 \\ & 1 & 1 & 2 \end{pmatrix}$
(a tri. with one extra diag.)

$$|H_n| = |H_{n-1}| + |H_{n-2}| \quad (\text{Fibonacci } F_{n+2})$$

• Markov matrix largest eigenvalue $\lambda_m = 1$. this eigenvc is the steady state. (A^∞)

Mar. mt. each column adds up to 1. (possibility transfer')

(proof: since $\lambda_A = \lambda_{A^T}$, thus λ can be 1; (proof)

and Gershgorin circle tells $\lambda - a_{ii} \leq \sum_{j \neq i} a_{ji}$, $\lambda \leq 1$)

6. Eigenvalues and Eigenvectors

6.1 introduction to eigenvalues

An eigen vector \vec{x} lies along the same line as Ax : $Ax = \lambda x$. The eigenvalue is λ .

Then $(A - \lambda I)x = 0$ and $A - \lambda I$ is singular and $\det(A - \lambda I) = 0$. So 1. Compute the det of $A - \lambda I$.

2. Find the n roots of this polynomial of degree n . 3. For each λ , solve $(A - \lambda I)x = 0$ to find an eigenvector x .

The eigenvalues of A^2 and A^{-1} are λ^2 and λ^{-1} , with the same eigenvectors.

(Vieta's formula or geometry) The sum of the λ 's equals the sum down the main diag of A or called the trace. $\sum_{i=1}^n \lambda_i = a_{11} + a_{22} + \dots + a_{nn} = \text{tr}(A)$. The products of the λ 's equals the det of A .

Special mts. (properties) lead to special eigenvalues and eigenvectors. $\prod_{i=1}^n \lambda_i = \det A$.

Proj. P have $\lambda = 1$ and 0 . Reflection R have $\lambda = 1$ and -1 . Rotation have $\lambda = e^{i\theta}$ and $e^{-i\theta}$.

Singular mt. has $\lambda = 0$. Triangular mt. has λ 's on its diagonal. Identity mt. has all n $\lambda = 1$.

Sym. mt. can be compared to a real number, antisym. mt. can be compared to an imaginary number.

Orthogonal mt. corresponds to a complex number with $|\lambda| = 1$. For these three special mts, their eigenvectors are perpendicular (conjugated) orthonormal.

6.2 diagonalizing a matrix

$Ax_i = \lambda x_i$ are the cols of $AX = X\Lambda$. The eigenvalue matrix Λ is diagonal and correspond with the order of x_i in X .

n ind. eigenvectors diagonalize A : $A = X\Lambda X^{-1}$. $\Lambda = X^{-1}AX$ (also diagonalize all powers $A^k = X\Lambda^k X^{-1}$)

No equal eigenvalues: ind. x from diff. λ . so eigenvectors are ind. X is inv. and A is diagonalizable.

Equal eigenvalues: A might have too few ind. eigenvectors. X^{-1} fails.

we have $GM \leq AM$ (geometric multiplicity \leq algebraic multiplicity)

when $GM < AM$, A is not diagg. $\dim(N(A - \lambda I))$ roots of $\det(A - \lambda I) = 0$

Similar matrices:

all $A = BCB^{-1}$. $A \sim C$. similar I .

they all share the eigenvalues of C .

Classical examples:

Markov $\begin{pmatrix} .8 & .3 \\ .2 & .7 \end{pmatrix} \Rightarrow \begin{pmatrix} .6 & .6 \\ .4 & .4 \end{pmatrix} \oplus I$

$\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ (Jordan form) $(\lambda = 1)$

$A^2 = A$. $\frac{xy^T}{y^T x}$ $(\lambda = 1.0)$

Fibnums. $\vec{u}_{k+1} = A\vec{u}_k$. $(u_k = A^k u_0)$ (difference eq.)

proof: multiply A several times to $\rightarrow U_j: c_1 x_1 + \dots + c_j x_j = 0 \Rightarrow \prod (c_i - \lambda_j)$ so $c_i = 0$. $c_i x_i = 0$ all other $c_j = 0$ ind. compressed in one pivot.

(ano2. way: $Sx = \lambda x, Sy = \beta y$. its colspace $C \perp$ nullspace N , so $x \perp y$ ($\lambda_1 \neq \lambda_2$); $S - \beta I$ is sym.

o' we can call R is S (sym.)'s sq root if $S = R^T R$. so from diaglize we

find R can be $= \Lambda^{\frac{1}{2}} Q$ (not unique, R to

对称阵性质证明
a proof of sym. mts has ortho. eigenvcs in algebraic way:

$$S = X \Lambda X^{-1}, S^T = X^{-T} \Lambda X^T \Rightarrow X^T X \Lambda = \Lambda X^T X \text{ or } \Lambda X^T X \text{ is sym.}$$

we know if A, B sym. (AB) sym. needs $AB = BA$. so Λ and $X^T X$ commutable and have same eigenvalues since Λ is diag, $X^T X$ must be like I or Λ , it's Q . (ano. way: $\lambda, x^T y = (Sx)^T y = x^T S y = \lambda_2 x^T y$

above has a premise that sym. mts can be diagonalized: ($x: \lambda_1, y: \lambda_2$) so $x^T y = 0$ ($\lambda_1 \neq \lambda_2$)

Schur's Theorem: Every sq. A factors into $Q T Q^{-1}$ where T is upper tri. and $Q^{-T} = Q^{-1}$. If A has real eivalues then Q and T can be chosen real: $Q^T Q = I$.

(informal: if S has repeated λ 's, slightly change S in diag $(c, 2c, \dots, nc)$. with $c \rightarrow 0$, orthonormal eivcs remained.)

real eigenvalues: $Sx = \lambda x \rightarrow S \bar{x} = \bar{\lambda} \bar{x}, \bar{x}^T S = \bar{x}^T \bar{\lambda}$.

$$x^T S x = x^T \lambda x \quad \bar{x}^T S x = \bar{x}^T \bar{\lambda} x \quad \Delta |A_{\pi\pi i}| = \text{positive or zero.}$$

Call mts have $\Pi d = \Pi \lambda$
same signs with pivots: $\bar{x}^T x$ is the squared length > 0 . (for antisym., $\lambda = -\bar{\lambda}$, it's pure imag.)

tri. fac. $S = LDL^T = \begin{pmatrix} 1 & & \\ & d & \\ & & 1 \end{pmatrix} \begin{pmatrix} d & & \\ & \dots & \\ & & 1 \end{pmatrix} \begin{pmatrix} 1 & & \\ & d & \\ & & 1 \end{pmatrix}$ so $\lambda = \bar{\lambda}$, it's real

when L moves to I ($0 \rightarrow 0$), $|D| I^T$ has eigenvalues = pivots d 's. (eivalue \rightarrow pivot) but we always have pivots d 's. \downarrow

But to change sign, a real eivalue would cross zero \rightarrow matrix singular at a moment X so no sign change in moving

幂零阵 nilpotent matrix: the following propositions are equivalent: ① A is nilpotent ($\exists m, s.t. A^m = 0$) ② $A^n = 0 \leftarrow A^{r(A)+1} = 0$ ③ all λ_i of $A = 0$

④ (from ③) $\det(A + kI) = k^n, \det(kA + I) = 1$ ⑤ $\forall k, \text{tr}(A^k) = 0$ (Cayley-Hamilton's Theorem: eigenpolynomial: $\det(\lambda I - A) = f(\lambda)$)

幂阵 $A^2 = I$: nothing new. $\lambda = \pm 1$. esp. sym. + ortho.) for nilp. A it's λ^n , then we have $f(A) = 0$.

更多关于矩阵指数: e^{At_1} times e^{At_2} must be $e^{A(t_1+t_2)}$, but $e^A e^B, e^B e^A, e^{A+B}$ can be all diff.

when A is tri., X, X^{-1}, e^{At} are also tri. (since $A - \lambda I$: $\begin{pmatrix} \lambda & & \\ & \lambda & \\ & & \lambda \end{pmatrix}$, X is tri.) $(e^A e^A = e^{2A})$

repeated λ : diaglize is impossible so compute e^{At} directly: some mt.'s series stops early.

$$A = \begin{pmatrix} 0 & 1 \\ -1 & 2 \end{pmatrix}, \lambda = 1, 1. \quad e^{At} = e^{It} e^{(A-I)t} = e^t (I + (A-I)t) \leftarrow (A-I)^2 \text{ is zero. (v.k.)}$$

antisym. $A = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$, then $e^{At} = \begin{pmatrix} \cos t & \sin t \\ -\sin t & \cos t \end{pmatrix}$, its eigenvalues are e^{it}, e^{-it} . so we find $t e^t$ here.

extension: Cosine form: solve $\frac{d^2 u}{dt^2} = -A^2 u$ $\cos At = I - \frac{1}{2!} A^2 t^2 + \frac{1}{4!} A^4 t^4 - \dots$ $(\cos At)^2 = -A^2 \cos At$ which corresponds.

Short form $u_{\cos} = \cos(At) u_{\cos}$ / specific form $u_{\cos} = c_1 x_1 + c_2 x_2$, $u_{\sin} = c_1 \cos \lambda_1 t x_1 + c_2 \cos \lambda_2 t x_2$

6.3 systems of differential equations

1) $\frac{du}{dt} = Au$. if $Ax = \lambda x$ then $u_{(t)} = e^{\lambda t} x$ will solve it. each λ and x gives a ind. solution.

2) start from $u_{(0)}$. it can be written as comb $c_1 x_1 + \dots + c_n x_n$ or $Xc = u_{(0)}$, $c = X^{-1} u_{(0)}$.

then each eigenvector is multiplied by $e^{\lambda_i t}$, the solution is $u_{(t)} = c_1 e^{\lambda_1 t} x_1 + \dots + c_n e^{\lambda_n t} x_n$.

3) (at past $\frac{d}{dt} \begin{pmatrix} y \\ z \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix}$ you may use $\frac{d}{dt} (y \pm z) = \pm (y \pm z)$ instead of $u_{(t)} = e^t \begin{pmatrix} 1 \\ 1 \end{pmatrix}$, $u_{(t)} = e^{-t} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$.)

Matrix exponential form: $e^{At} = I + At + \frac{1}{2!} (At)^2 + \dots$, $(e^{At})' = A e^{At}$.

2. e^{At} 's eigenvalues are $e^{\lambda t}$ (sub $A = X \Lambda X^{-1}$ into exponential, 3. e^{At} has inverse e^{-At} , when A is

so $u_{(t)} = \sum_i c_i e^{\lambda_i t} x_i = X e^{At} X^{-1} u_{(0)} = e^{At} u_{(0)}$. antisym, $(A^T = -A)$ e^{At} is orthogonal.

(we call A is stable when $e^{At} \rightarrow 0$ thus all eigenvalues of A have real part < 0 , imaginary part

(Not included: if two λ 's are equal, with only one eigenvector, another solution is needed: $t e^{\lambda t} x$)

6.4 symmetric matrices

Every sym. matrix S can be diagonalized, and has

(Principle Axis / Spectral Theorem) $S = Q \Lambda Q^T$ n real eigenvalues λ_i and n orthonormal eigenvectors q_1, \dots, q_n .

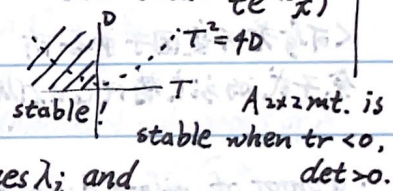
The number of positive eigenvalues of S = The number of positive pivots of S (proof via q_i 's)

(Antisym. matrix $A^T = -A$ have imag. λ 's and ortho. complex 'signs matched' q 's)

For real matrices (nonsym.), complex λ 's and x 's come in conjugate pairs: $Ax = \lambda x$

(every ortho. Q have $| \lambda | = 1$ since then)

$$A \bar{x} = \bar{\lambda} \bar{x}$$



6.5 positive definite matrices

positive definite (sym.) S : all eigenvalues $> 0 \iff$ all pivots $> 0 \iff$ all upper left det > 0 .

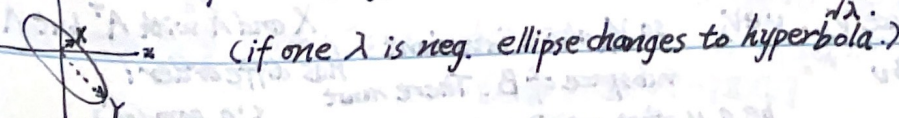
*('Energy-based def/test') $x^T S x > 0$ for all vecs $x \neq 0$. (so obviously S pos. T pos. $\rightarrow S + T$ pos.)

5. A col ind. (to rule out zero), $S = A^T A$, $(A = chol(S) = L \sqrt{D} / A = Q \Lambda Q^T \text{ etc.})$ ($x^T S x = \|Ax\|^2 > 0$)

positive semidefinite: allows $\lambda = 0$, pivot = 0, det = 0, $x^T S x = 0$ ($x \in N(S)$), A has dep. cols.

geometry: The graph of $x^T S x = 1$ is an ellipse: $(x \ y) Q \Lambda Q^T \begin{pmatrix} x \\ y \end{pmatrix} = (X \ Y) \Lambda \begin{pmatrix} X \\ Y \end{pmatrix} = \lambda_1 X^2 + \lambda_2 Y^2 = 1$. (standard form)

The axes point along eivectors of S , half-lengths are $\frac{1}{\sqrt{\lambda_i}}$ ($X = q_1^T x$, $Y = q_2^T x$)



o' S pos. T pos. \rightarrow ST might not sym

but pos. eivalues: $STx = \lambda x$.

$$\lambda = \frac{(Tx)^T STx}{\lambda x^T Tx} > 0$$

(Normal includes sym. antisym. ortho.)

o N has n orthonormal eigenvectors ($N = Q\Lambda Q^H$) if and only if N is normal. ($N^H N = N N^H$)

(proof: o $N = Q\Lambda Q^H$ subin. @ Schur T: $N = Q\Lambda Q^H$ and T is Λ .)

o' sym. S: o no diag entry can be larger than λ_{max} (use $S = \lambda_i p_i p_i^T$)

@ diag entries fall in between the λ 's. (S_{xx} : $\frac{a_{ii}}{\det(A - \lambda I)}$)

□ Congruence . similarity — same transform on diff. basis.

合同矩阵 congruence — same bilinear / (co,2) tensor / metric on diff. basis.

$$x^T A x$$

$$T = T_{ij} e_i e_j^T, \text{ if } e_i = R_{ij} e_j' \text{ then } T'_{ij} = T_{kl} R_{ki} R_{jl}, \text{ mt. form: } T' = R^T T R.$$

$$x^T \underbrace{C^T A C}_B x$$

some quadratic surface on diff. basis: $\textcircled{+} * \textcircled{+}$

. In real, congruence \leftrightarrow equal inertia index. keep the pos./neg. inertia index.

In complex, congruence \leftrightarrow equal rank.

□ Jordan canonical form

约当标准形 (规范) . sq. $A \sim J = \begin{pmatrix} J_1 & & \\ & J_2 & \\ & & \ddots \\ & & & J_m \end{pmatrix}$, $J_i = \begin{pmatrix} \lambda_i & & \\ & \ddots & \\ & & \lambda_i \end{pmatrix}$

<可参考 J_i^k : Jordan pot.>

<可参考不变因子中各阶 $g_m =$ number of Jordan blocks. $\leq m =$ scale/cols of Jordan blocks.

余子式的方法替代 (M.A.M.)

$$A = X J X^{-1}. AX = X J \text{ solve } x_i \text{ one by one. } \begin{pmatrix} x_1, x_2, x_3, \dots \\ \lambda_i \\ \vdots \end{pmatrix} \begin{matrix} \text{call } \lambda_i \text{ adds} \\ (A - \lambda_i I) x_1 = 0 \\ (A - \lambda_i I) x_2 = x_1 \\ \vdots \end{matrix}$$

. a proof of nilpot. A has $A^{r(A)+1} = 0$.

' a more classical way of Jordan: $(A - \lambda I) x_j = x_{j-1}$

since $\lambda = 0$, $A \sim J = \begin{pmatrix} J_{(c_1, k_1)} & & \\ & \ddots & \\ & & J_{(c_m, k_m)} \end{pmatrix}$

@ Primary Decomp. $f(\lambda) = (\lambda - \lambda_1)^{m_1} \dots (\lambda - \lambda_m)^{m_m}$

$r(A) = r(J) \geq r(J_{(c_i, k_i)}) = k_i - 1$ so $k_i \leq r(A) + 1$. $V_i = \ker(A - \lambda_i I)^{k_i}$ then $V = \bigoplus_{i=1}^m V_i$ (根子空间)

$$J_{(c_i, k_i)}^{k_i} = 0 \rightarrow A^{r(A)+1} = 0$$

@ Cyclic Decomp. $V_i = \bigoplus_{j=1}^{t_i} \text{span}\{N^j \omega_i : 0 \leq j < t_i\}$

<可参考 Young 图> ($t = \dim(\ker N)$, $N = A - \lambda I$) (强循环子空间)

o rank-one decomp. / Spectral decomp.: a synthesis concept of decomp. $A = \sum_i \lambda_i G_i = \sum_i \lambda_i u_i v_i^T$

秩一分解

谱分解

($G_i G_j = \delta_{ij} G_i$ if S) G_i is idempot. λ_i is the eigenvalue. ($G_i = p_i p_i^T$ or $u_i v_i^T$)

I generalize it: all n rank-one mts $u_i v_i^T$ adds up to A. Obviously it's not unique.

△ a rank r mt. can be expanded by r rank-one mts and at least r. ($r = r - k + k$, general)

proof: A can be elementary transformed to $P \begin{pmatrix} I_r & \\ & 0 \end{pmatrix} Q$ so $A = \sum_{i=1}^r p_i \cdot 1 \cdot q_i^T$

and since $r(A+B) \leq r(A) + r(B)$, at least should r rank-one mts.

□ Table of Eigens [Matrix/ λ / x] P362

特殊特征值向量表

△ \mathbb{Q} has all singular values = 1 as well as I.

o a proof of commutable mts have common eigenvcs: esp. normal mt A have same

A has a λ and its eigensubspace V_λ . $v \in V_\lambda$, so V_λ also a invariant X and Λ with A^T but Λ

$$ABv = BA v = \lambda Bv$$

subspace of B. There must has diff. order!

be a μ that s.t. $Bu = \mu u$ with $Au = \lambda u$. ('a paradox')

Then u is the common eivec.

7. The Singular Value Decomposition (SVD)

7.1 image processing by linear algebra

The SVD separates any matrix A into rank-one pieces $A = \sum_i \sigma_i u_i v_i^T$. in order of importance $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r > 0$.

The choices of u 's and v 's are orthon. $u_i^T u_j = \delta_{ij}$, $v_i^T v_j = \delta_{ij}$.

\triangle has infinite rank, $\equiv \square$ has low finite rank. Discard small σ 's. (the dogma of image process.)

7.2 bases and matrices in the SVD $A = U \Sigma V^T = u_1 \sigma_1 v_1^T + \dots + u_r \sigma_r v_r^T$

one way to deduce: $A^T A v_i = \sigma_i^2 v_i \rightarrow A v_i = \sigma_i u_i \rightarrow A^T u_i = \sigma_i v_i \rightarrow A A^T u_i = \sigma_i^2 u_i$
 1. both $A^T A$, $A A^T$ sym. pos. 2. choose $u_i \in \mathcal{C}(A)$. 3. dual to 2. 4. u_i are eigenvectors of $A A^T$. so again $u_i \perp u_j$.
 v_i chosen to be eigenvectors of $A^T A$. ($u_1 \dots u_r$) and we $u_i \in \mathcal{R}(A^T) = \mathcal{C}(A)$ and left $u_{r+1} \dots u_m \in \mathcal{N}(A^T)$.
 with $\sigma_i > 0$. (r nonzero) can proof: $u_i^T u_j = \frac{v_i^T A^T A v_j}{\sigma_i \sigma_j} = \frac{\sigma_j}{\sigma_i} v_i^T v_j = 0$.
 so $v_i \perp v_j$. $v_1 \dots v_r \in \mathcal{R}(A)$ and $u_i^T u_i = v_i^T v_i = 1$ also.

To conclude: $\sigma_1^2 \dots \sigma_r^2$ are the nonzero eigenvalues of $A^T A$ and $A A^T$.

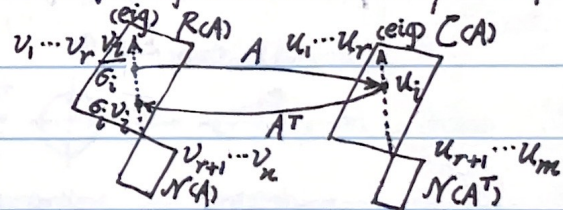
the orthon. cols of U and V are the eigenvectors of $A^T A$ and $A A^T$.

Also those cols hold orthon. bases for the four fundamental subspaces of A .

then diagonalize $A: AV = U \Sigma$.

(reduced SVD: $(u_1 \dots u_r) \begin{pmatrix} \sigma_1 & & \\ & \dots & \\ & & \sigma_r \end{pmatrix} \begin{pmatrix} v_1^T \\ \vdots \\ v_r^T \end{pmatrix}$;
 $m \times r$ $r \times r$ $r \times n$)

full SVD: $(u_1 \dots u_m) \begin{pmatrix} \sigma_1 & & \\ & \dots & \\ & & \sigma_r \end{pmatrix} \begin{pmatrix} v_1^T \\ \vdots \\ v_n^T \end{pmatrix}$.
 $m \times m$ $m \times n$ $n \times n$)



esp. When A is pos. (semipos.) sym. mat and only when

can $A = X \Lambda X^{-1} = Q \Lambda Q^T = U \Sigma V^T$ ($X=U=V$; $\lambda \geq 0, \Lambda = \Sigma$)

Singular versus Eigen:

$$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 2 \\ 0 & 0 & 3 \\ 1 & 0 & 0 \\ \hline 0 & 0 & 0 \end{pmatrix}$$

$\lambda = 0, 0, 0, 0$ jump to $\frac{\pm 1}{10}, \frac{\pm i}{10}$.

This shows serious instability of eigenvalues when $A A^T$ is far from $A^T A$.

The singular values are always stable. σ_4 emerges as $\frac{1}{60000}$ from 0. $\sigma = 3.2.1.0$

(At the other extreme, $A A^T = A^T A$ 'normal' means A has orthon. eigenvecs and stable eigenvalues.)

since $\sigma = \sqrt{\lambda}$, we have (for sq.)

one way one-at-a-time instead of all-at-once:

$S = A^T A = (V \Sigma^T \Sigma V^T)$ $\prod \lambda_i = \prod \sigma_i^2 = \det A$
 $S = \lambda_1 u_1 u_1^T + \dots + \lambda_r u_r u_r^T$; $A = \sigma_1 u_1 v_1^T + \dots + \sigma_r u_r v_r^T$

$\lambda_1 = \max \frac{x^T S x}{x^T x}$ when $x = q_1 = v_1$
 $\sigma_1 = \max \frac{\|A x\|}{\|x\|}$ when $x = v_1$ ($\max \frac{x^T A^T A x}{x^T x}$)
 $\lambda_2 = \max \frac{x^T S x}{x^T x}$ with $q_1^T x = 0$ when $x = q_2 = v_2$
 $\sigma_2 = \max \frac{\|A x\|}{\|x\|}$ with $v_1^T x = 0$ when $x = v_2$

(Rayleigh quotient $r(x) = \frac{x^T S x}{x^T x}$. you can see $Q^T x$ as x . the $x^T Q^T \Lambda Q x = \sum \lambda_i \cos^2 \theta_i$. choose λ .
 2. $\frac{\partial r}{\partial x} = 0$. For next step, $Q^T x$ has first component = 0

or $Q_1^T S Q_1 = \begin{pmatrix} \lambda_1 & \\ & S_{n-1} \end{pmatrix}$

$\|A x\| = \|U \Sigma V^T x\| = \|\Sigma V^T x\| \leq \sigma_1 \|V^T x\| = \sigma_1 \|x\|$

o' 补充说明: when q_1, \dots, q_n ortho., we have $q_1 q_1^T + \dots + q_n q_n^T = I$. vice versa. (?)

o $Q_1 A Q_2^T$ has the same σ 's as A and $Q S Q^T$ has the same λ 's as S .

Computing of singular: $Q^T S Q \rightarrow$ triang. $Q_1 A Q_2^T \rightarrow$ bidiag.

o $\sigma_{\max} \geq |\lambda|_{\max}$; $\sigma_{\min} \leq |\lambda|_{\min}$. (以前) $\|Ax\| \leq \sigma \|x\|$; $\|Ax\| \geq \sigma_{\min} \|x\|$ or $\frac{1}{\sigma_{\min}} \geq \left| \frac{1}{\lambda_{\min}} \right|$
 = $\lambda \|x\|$ inverse to A^{-1}

□ Norm of matrices: 向量 p 范数 $\|x\|_p = (\sum_j |x_j|^p)^{1/p}$ ($1 \leq p$) 范数满足 ① $\|x\| \geq 0$ ② $\|\alpha x\| = |\alpha| \|x\|$ ③ $\|x+y\| \leq \|x\| + \|y\|$.
 矩阵范数 范数满足 ① $\|x\| \geq 0$ ② $\|\alpha x\| = |\alpha| \|x\|$ ③ $\|x+y\| \leq \|x\| + \|y\|$.
 常 λ 1-norm, 2-norm, ∞ -norm (Euclidean) $(\max_j |x_j|)$

inducing norm/operator norm: $\|A\|_p = \max_{x \neq 0} \frac{\|Ax\|_p}{\|x\|_p}$. 常 λ $\|A\|_1 = \max_j \sum_i |a_{ij}|$ (max col sum)

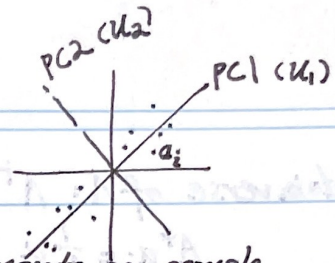
Frobenius norm (F-norm): $\|A\|_F = \sqrt{\sum_{i,j} |a_{ij}|^2} = \sqrt{\text{tr}(A^T A)}$ $\|A\|_2 = \sqrt{\lambda_{\max}(A^T A)} = \sigma_1$ (Enc./spectron)

theorem 1: $\rho(A) \leq \|A\|$ (ρ is spectrum radius, $\rho(A) = \max |\lambda_i|$)
 $(\forall \epsilon > 0, \exists \|\cdot\|_p$ st. $\|A\|_p \leq \rho(A) + \epsilon)$
 eq. when $\|\cdot\|$ is 2-norm and A is sym.

theorem 2 (Gelfand formula):
 $\rho(A) = \lim_{k \rightarrow \infty} \|A^k\|^{1/k}$

o **rank-one mt.** $A = \sigma u v^T$ (unit vec u and v), $A^+ = \frac{v u^T}{\sigma}$. $AA^+ = u u^T$, $A^+A = v v^T$.

o' SVD of A is equivalent to the diagonalization of sym. $M = \begin{pmatrix} 0 & A^T \\ A & 0 \end{pmatrix}$, with eivec = $\begin{pmatrix} v \\ u \end{pmatrix}$.



7.3 principal component analysis (PCA by the SVD)

Data comes in a matrix: n samples and m measurements per sample.

Subtracting the mean of each row to center A . Then SVD finds the (\vec{u}, \vec{v})

(sample) covariance matrix: $S = \frac{AA^T}{n-1}$ 'eigen-' or 'comb-' that contains the most info. or explains the most variance.

Then we find λ and σ , u and the leading direction in the scatter plot.

Total Variance $T = s_1^2 + \dots + s_m^2 = \sigma_1^2 + \dots + \sigma_m^2 = \text{tr}(S)$. PC1 accounts for $\frac{\sigma_1^2}{T}$ of total var.

PCA = perpendicular least squares (orthogonal regression): (not OLS / vertical least squares!)

$$\sum_{j=1}^n \|a_j\|^2 = \|A^T u_1\|^2 + \|A^T u_2\|^2 = \sum_{j=1}^n |a_j^T u_1|^2 + \sum_{j=1}^n |a_j^T u_2|^2$$

$A^T A x = A^T b$
When $u_1^T A A^T u_1$ maximized, the squared distances sum minimize.

Correlation Matrix: $C = \frac{DA^T A D}{n-1} = \text{DSD}$. (and u_2 the second direction in ortho) measurements may have diff units and original scaling is not meaningful.

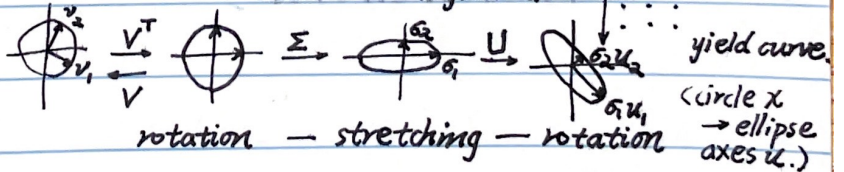
DA makes row have same length $\sqrt{n-1}$. or makes S has 1's on diag by dividing s_1, \dots, s_n .

App. Examples: \odot Genetic variation in Europe. \otimes Eigenfaces. \odot Model order reduction, POD.

(Δ When A is centered and $m > n$, σ 's have at most $n-1$ nonzeros.)

\oplus Web matrix (Pagerank) \odot $T^t \dots$ iter.

7.4 the geometry of the SVD



The norm $\|A\| = \max \frac{\|Ax\|}{\|x\|} = \sigma_1$.

tri.ineq. $\|A+B\| \leq \|A\| + \|B\|$, product ineq. $\|AB\| \leq \|A\| \|B\|$ (proof: $\|Ax\| \leq \|A\| \|x\|$, times x)

Eckart-Young (Mirsky) Theorem: The closest rank k mt to A is $A_k = \sigma_1 u_1 v_1^T + \dots + \sigma_k u_k v_k^T$. ($\|A-B\| \geq \|A-A_k\| = \sigma_{k+1}$, $\forall B$ of $\text{rc}(B) = k$.)

Polar Decomp.: like $e^{i\theta} r$, $A = QS$ where Q is orthogonal and S is sym. pos. semidefinite. rotation · stretching $Q = UV^T$, $S = V\Sigma V^T$. if A is inv., S is pos. definite.

Δ $Q = UV^T$ is the nearest ortho. mt to A . (cor S is the sq. root of $A^T A = V\Sigma^2 V^T$.) (also $A = kQ$ with $k = U\Sigma U^T$ and Q the same) ($\|Q - A\|_{\text{min}}$.)

A_0 is the nearest sin. mt to A by changing σ_{min} to zero.

$n \times m$

pseudoinverse of A : $A^+ = V \Sigma^+ U^T$ $(v_1 \dots v_n) \begin{pmatrix} \sigma_1^{-1} & & \\ & \ddots & \\ & & \sigma_r^{-1} \end{pmatrix} (u_1 \dots u_m)^T$

$A^+ u_i = \frac{1}{\sigma_i} v_i$ ($i \leq r$) and $A^+ u_i = 0$ ($i > r$) $n \times n$ $n \times m$ $m \times m$ always exists.

$N(A^+)^T = N(A)$
 $N(A^+) = N(A^T)$
 $C(A^+) = R(A)$
 $R(A^+) = C(A)$

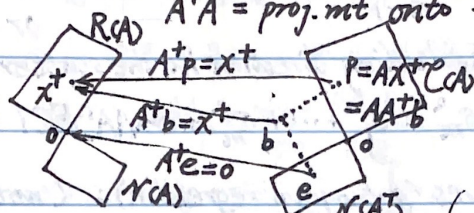
If A^{-1} exists, A^+ is the same as A^{-1} in that case $m=n=r$ and $V \Sigma^{-1} U^T$.

since $\Sigma^+ \Sigma = \begin{pmatrix} I & 0 \end{pmatrix}$, we get

$AA^+ = A^+A = I$
 (proj. itself)

$AA^+ = \text{proj. onto the col space of } A.$

$A^+A = \text{proj. onto the row space of } A.$



A^+ can get Ax^+ in col space back to $A^+Ax^+ = x^+$ in row space

$A^+A = \begin{pmatrix} I & 0 \\ 0 & 0 \end{pmatrix}$ $n \times n$ row space $= \sum v_i v_i^T$
 nullspace $= \text{an ex.}$
 $AA^+ = \begin{pmatrix} I & 0 \\ 0 & 0 \end{pmatrix}$ $m \times m$ col space $= \sum u_i u_i^T$
 left nullspace $=$

least squares with dependent cols: $A^T A x = A^T b$ has inf. many solutions when A is singular.
 pseudoinv. gives a way to choose a best solution: $x^+ = A^+ b$ which is the shortest

$(\|x^+\|_{\min}, \text{ no nullspace part})$ solution to $A^T A \hat{x} = A^T b$ and $A \hat{x} = p$.

If A is full col rank ($r=n$) then it has left inv. $L = (A^T A)^{-1} A^T$, $LA = I$. $A^+ = L$ in this case;

If A is full row rank ($r=m$) then it has right inv. $R = A^T (A A^T)^{-1}$, $AR = I$. $A^+ = R$ in this case.

(proof: LA and AL are proj's onto row space (\mathbb{R}^n) and col space (sub of \mathbb{R}^m , the origin P);
 AR and RA the same.)

8. Linear Transformations

8.1 the idea of a linear transformation

A transformation assigns an output $T(v)$ to each input vector $v \in V$.

Linearity requires $T(cv+dw) = cT(v) + dT(w)$. Note $T(0) = 0$, so affine trans. isn't linear.

Lines \rightarrow Lines, Triangles \rightarrow Triangles, Comb \rightarrow Comb, so basis is important. $(T(v) = Av + u_0, \text{ ex.})$

$A = \begin{pmatrix} 0 & 1 & & 0 \\ & 0 & 2 & \\ & & 0 & \dots \\ & & & 0 \end{pmatrix}$ = mt. form of the derivative $T = \frac{d}{dx}$, with basis $v_i = x^{i-1}$, nullspace a line in func. sp.

and integral give the pseudoinv. $T^+ = \int_0^x dx$, (mt. form $A^+ = \begin{pmatrix} 0 & \dots & 0 \\ 1 & & 0 \\ & \ddots & \\ 0 & & \frac{1}{n} \end{pmatrix}$) $A^+A = (0 \ I)$, $AA^+ = I$.

The product of two trans. $(ST)(v) = S(T(v))$ still linear. Inverse trans. T^{-1} brings every vec

$T: V \rightarrow W$ Range of T = set of outputs. Kernel of T = set of all inputs for which $T(v) = 0$.
 Ex. $T(v) = Av$ (Image) $\mathcal{C}(A)$ $\mathcal{N}(A)$ $T(v)$ back v .

8.2 the matrix of a linear transformation

Every lin. trans. from V to W can be converted to a matrix. This mt. depends on bases.

Cols of 1 to n of the matrix will contain those outputs $T(v_1)$ to $T(v_n)$.

$T(v_i) = \text{comb of output basis vectors} = a_{i1}w_1 + \dots + a_{in}w_n$.

For every $v \in V$ that $v = c_1v_1 + \dots + c_nv_n$, linearity gives $T(v) = c_1T(v_1) + \dots + c_nT(v_n)$, as Ac .

Multiplication corresponds. $TS: U \rightarrow V \rightarrow W$, $AB: (m \times n)(n \times p) = (m \times p)$.
 $\mathbb{R}^p \quad \mathbb{R}^n \quad \mathbb{R}^m$

Change of basis: $Vc = V'c'$, (same vec written in diff basis)
 1. $V'B = V, B = V^{-1}V; C' = Bc$.
 2. $A' = B_{out}^{-1}AB_{in}$ (B contains new basis vecs b in cols written in the standard/old basis.)
 (of course it's inv.)

8.3 the search for a good basis

vec. sp. 1. $B_{in} = B_{out} = \text{eigenvec mt. } X, X^{-1}AX = \text{eigenvalues in } \Lambda$. (sq. A and has n ind. eivcs.)

2. $B_{in} = V$ and $B_{out} = U$: singular vecs of $A, U^{-1}AV = \text{diag } \Sigma$. — 'isometric' (def: $c = \alpha_1^T A \alpha_2$ when α_1, α_2 are ortho.)

3. $B_{in} = B_{out} = \text{generalized eigenvecs of } A, B^{-1}AB = \text{Jordan form } J$. (sq. A but only $s \leq n$ ind. vecs.)

$J = \begin{pmatrix} J_1 & & \\ & \ddots & \\ & & J_s \end{pmatrix}$, Jordan block: $J_i = \begin{pmatrix} \lambda_i & 1 & & \\ & \lambda_i & \ddots & \\ & & \ddots & 1 \\ & & & \lambda_i \end{pmatrix}$ $(A - \lambda_i I)b_{j+1} = b_j$ (at one block)
 $(J - \lambda_i I)x_{i+1} = x_i$ instead of 0, $x_j = (c_0 \dots c_{j-1}, 0 \dots 0)$.

Matrices are similar if they share the same Jordan form, not otherwise. and $Bx_j = b_j$.

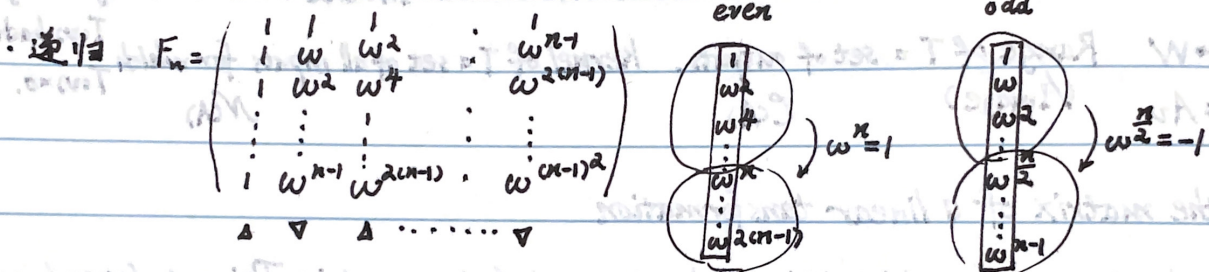
In app. we take J^k and e^{Jt} involves $e^{\lambda t}$ times powers $1, t, \dots, t^{s-1}$.

(Jordan form is unstable!)

o 矩阵收敛性(四):

Hilbert matrix $H = \begin{pmatrix} 1 & \frac{1}{2} & \frac{1}{3} & \frac{1}{4} \\ \frac{1}{2} & \frac{1}{3} & \frac{1}{4} & \frac{1}{5} \\ \frac{1}{3} & \frac{1}{4} & \frac{1}{5} & \frac{1}{6} \\ \frac{1}{4} & \frac{1}{5} & \frac{1}{6} & \frac{1}{7} \end{pmatrix}$, $H_{ij} = \frac{1}{i+j-1}$ if use powers $1, x, x^2, x^3 \dots$ as a func basis.
 $B^T B$ is not ortho. and equals H : bad!
 H is ill-conditioned.

o FFT 形为化补充:



o proof of Schur Theorem: $U^H A U = T$.
 (Schur decomp.) ($A \in \mathbb{C}^n$, U unitary, T tri.)

induction: find A 's λ_1 a eigenvector ϵ_1 , expand it to a basis.

$U_1 = (\epsilon_1, \eta_2, \dots, \eta_n)$ $U_1^H A U_1 = \begin{pmatrix} \lambda_1 & \dots \\ 0 & A_1 \end{pmatrix}$ then do the same thing on A_1 , till A_{n-1} .

finally $(U_1 U_2 \dots U_{n-1})^H A (U_1 U_2 \dots U_{n-1}) = \begin{pmatrix} \lambda_1 & \dots \\ \vdots & \ddots \\ 0 & \dots & \lambda_n \end{pmatrix}$ $U_2^H A_1 U_2 = \begin{pmatrix} \lambda_2 & \dots \\ 0 & A_2 \end{pmatrix}$, $U_2 = \begin{pmatrix} 1 & 0 \\ 0 & U_2 \end{pmatrix}$.

Schur inequality: $\sum_i |\lambda_i|^2 \leq \sum_i \sum_j |a_{ij}|^2$ (proof: $\sum_i |\lambda_i|^2 = \sum_i |t_{ii}|^2 \leq \sum_i |t_{ii}|^2 + \sum_{i \neq j} |t_{ij}|^2 = \text{tr}(T T^H)$)

o change in A^{-1} from a change in A : Sherman - Woodbury - Morrison formula: $A^{-1} + A^{-1} U (C - V^T A^{-1} U)^{-1} V^T A^{-1}$ (eg. when A is normal i.e. unitary similar diag) not inv. if $V^T U = I$.

- $M = I - uv^T$, $M^{-1} = I + \frac{uv^T}{1 - v^T u}$ (rank-one change) $\Delta Ax = b \rightarrow My = b$: ($M = A - uv^T$)
- $M = A - uv^T$, $M^{-1} = A^{-1} + \frac{A^{-1} uv^T A^{-1}}{1 - v^T A^{-1} u}$
- $M = I - UV$, $M^{-1} = I_n + U (I_m - VU)^{-1} V$
- (matrix inversion lemma) $M = A - UW^{-1}V$, $M^{-1} = A^{-1} + A^{-1} U (W - V A^{-1} U)^{-1} V A^{-1}$

o Hamilton-Cayley Theorem (见前) $f(\lambda) = \det(\lambda I - A)$ then we have $f(A) = 0$.

corollary: A^k can be simplified by $f(A) = 0$ (downgrade) diagonalizable mts. are dense.
 A^{-1} can be expanded by a gca ($f(A)$ times A^{-1} to get)

(DFT)

4. Bin = Bant = Fourier matrix F, Fx is a Discrete Fourier Transformation of x.

the eigenvector mt. F diagonalizes the permutation mt. P: ex. $\begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$ $\Lambda = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$ $F = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix}$

F is sym. ortho? (unitary.) Circulant matrix: $C = \begin{pmatrix} c_0 & c_1 & c_2 & c_3 \\ c_3 & c_0 & c_1 & c_2 \\ c_2 & c_3 & c_0 & c_1 \\ c_1 & c_2 & c_3 & c_0 \end{pmatrix}$ eigenvalues $c_0 + c_1\lambda + c_2\lambda^2 + c_3\lambda^3$
func.sp. its four eigenvalues are given by $Fc = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix} \begin{pmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \end{pmatrix} \rightarrow \begin{matrix} c_0(\lambda = \pm 1, \pm i) \\ \end{matrix}$

5. The Fourier basis. $1, \sin x, \cos x, \sin 2x, \cos 2x \dots$ periodic and orthogonal, so Fourier coeff.

6. The Legendre basis. $1, x, x^2 - \frac{1}{3}, x^3 - \frac{3}{5}x \dots$ ex. $a_i = \frac{\int f(x) \cos i x dx}{\int \cos i x \cos i x dx}$

come from applying the G-S idea to orthogonalize $1, x, x^2, x^3 \dots$ ex. $\frac{(x^3, x)}{(x, x)} = \frac{3}{5}$, so $x^3 - \frac{3}{5}x \perp 1, x, x^2$

7. The Chebyshev basis. $1, x, 2x^2 - 1, 4x^3 - 3x \dots$ $\cos n\theta \rightarrow \cos \theta$ (remind even \perp odd (x))

9. Complex Vectors and Matrices

9.1 complex numbers complex plane. polar form. n th roots of 1. $|z|^n = r^n = z\bar{z}$, $\omega = e^{\frac{2\pi i}{n}}$. $\frac{1}{z} = \frac{\bar{z}}{|z|^2}$, $\bar{\bar{z}} = z$.

9.2 hermitian and unitary matrices

Real versus Complex in a nutshell: When you transpose vec z or mt A , take the complex conjugate too.
 $\bar{z}^T = z^H$ called Hermitian or adjoint. (For truly useful $u^H u = \|u\|^2$ not $u^T u$.)

dot.p / inner.p: $u^H v$ the order is now important $u^H v \neq v^H u$. $(Au)^H v = u^H (A^H v)$
 $(AB)^H = B^H A^H$

$S = S^H$ called Hermitian matrix. $z^H S z$ is real for $\forall z$; all eigenvalues of S is real;

$Q^H Q = I$, $Q^H = Q^{-1}$ (sq.) called unitary. eigenvcs from diff eigenvalues are ortho. $q_i^H q_j = 0$

$\|Qz\| = \|z\|$ and then $Qz = \lambda z$ leads to $|\lambda| = 1$.

Δ good row space is no longer $\mathcal{C}(A^T)$ but $\mathcal{C}(A^H)$.

9.3 the Fast Fourier Transform (FFT) $F_n c = \begin{pmatrix} 1 & 1 & \dots & 1 \\ \omega & \omega^2 & \dots & \omega^{n-1} \\ \omega^2 & \omega^4 & \dots & \omega^{2(n-1)} \\ \vdots & \vdots & \ddots & \vdots \\ \omega^{n-1} & \omega^{2(n-1)} & \dots & \omega^{(n-1)^2} \end{pmatrix} \begin{pmatrix} c_0 \\ c_1 \\ c_2 \\ \vdots \\ c_{n-1} \end{pmatrix} = y = \begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_{n-1} \end{pmatrix}$ $\omega = e^{\frac{2\pi i}{n}}$

$$F_n^H = \bar{F}_n, F_n \bar{F}_n = nI, F_n^{-1} = \frac{1}{n} \bar{F}_n$$

frequency space \xleftrightarrow{F} physical space
 $\xleftarrow{F^{-1}}$

(with $w = \bar{\omega} = e^{-\frac{2\pi i}{n}}$, or permute row $i \leftrightarrow N-i$)

$\Delta y_j =$ the Fourier series $\sum_k c_k e^{ikx_j}$

those n points x (angles) equally spaced around.

Δ entry in row j col k is ω^{jk} , zeroth row and col contains all $\omega^0 = 1$.

Δ Fr.mt. is the Vandermonde mt. for interpolation at n .

recursion: $F_n \rightarrow F_{\frac{n}{2}}$

$$F_n = \begin{pmatrix} I_{\frac{n}{2}} & D_{\frac{n}{2}} \\ I_{\frac{n}{2}} & -D_{\frac{n}{2}} \end{pmatrix} \begin{pmatrix} F_{\frac{n}{2}} \\ F_{\frac{n}{2}} \end{pmatrix} \begin{pmatrix} 1 & \dots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \dots & 1 \end{pmatrix} \begin{matrix} \text{even-odd perm.} \\ y_j = \sum_0^{n/2-1} \omega^{jk} c_{2k} + \sum_0^{n/2-1} \omega^{j(2k+1)} c_{2k+1} \\ (\omega' = \omega^2) = \sum_0^{n/2-1} \omega'^{jk} c_{2k} + \sum_0^{n/2-1} \omega'^{jk} c_{2k+1} \cdot \omega^j = y'_{\text{even}} + \omega^j y'_{\text{odd}} \end{matrix}$$

$$D = \text{diag}(1, \omega, \dots, \omega^{\frac{n}{2}-1})$$

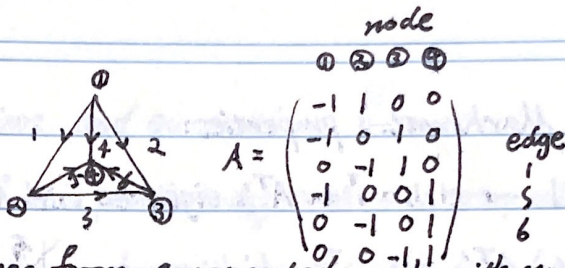
$$O(n^2) \rightarrow O\left(\frac{1}{2} n \log_2 n\right)$$

(C's final permutation:

$$\begin{matrix} 0 & 1 & 0 & 1 \\ \leftarrow & & & \end{matrix} \text{ bit-reversed order}$$

10. Applications

10.1 graphs and networks



The incidence matrix A comes from a connected graph with n nodes and m edges.

Elimination reduces every graph to a tree. Rows' dependency \leftrightarrow edges form a loop.

$N(A)$: The constant vec (c, c, \dots, c) make up the nullspace, $\dim=1$.

$C(A^T)$: The edges of any tree give r ind. rows, $r = n - 1$.

Δ Use orthogonality to judge a vec whether in a space of A .

$C(A)$: Ax gives voltage diffs. Voltage Law: The components of Ax add to zero around all loops, $\dim = n - 1$.

$N(A^T)$: Current Law: $A^T y = 0 = (\text{flow in}) - (\text{flow out})$ is solved by loop currents, since Ohm's Law $y = -CAx$. C is the conductance mat. (diag) there are $m - r = m - n + 1$ ind. (small) loops.

if all $c = 1$, we get the (graph) Laplacian matrix $A^T A$. (With A and C , Nodes - Edges + Loops = 1. we call it a Network.) Δ Euler's Formula \dim small

10.2 matrices in engineering

$A^T C A x = f$ (batteries / current sources \dots)
 $-\frac{d}{dx} (C(x) \frac{d}{dx}) = f(x)$ with boundary conditions: $\Delta A^T A$ is not inv. one node has to be grounded. ($x_j = 0$)

divide the bar into n pieces of Δx ; $u(0) = 0$ and $u(1) = 0$ or $u(1) = 0$ or $u(0) = 0$ or $u(1) = 0$ Thus remove one row and col.

replace $\frac{d}{dx}$ by A and $-\frac{d(\cdot)}{dx}$ by A^T ; (they include $\frac{1}{\Delta x}$) $\frac{du}{dx}(1) = 0$ ('fixed-free' or 'fixed-fixed')

end conditions are $u_0 = 0$ and $(u_n = 0 \text{ or } y_n = 0)$; $(Ku = f)$

$C(x)$ corresponds to C , thus: $f = A^T y$, $y = Ce$, $e = Au$ give $A^T C A u = f$.

we have three choices in A replacing $\frac{d}{dx}$: $\frac{u(x+\Delta x) - u(x)}{\Delta x}$ (forward), $\frac{u(x) - u(x-\Delta x)}{\Delta x}$ (backward), $\frac{u(x+\Delta x) - u(x-\Delta x)}{2\Delta x}$ (centered). Δ fix, free

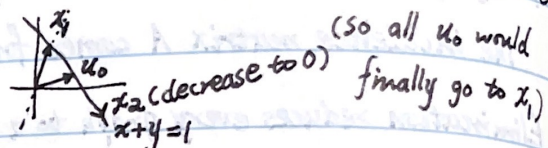
fixed-fixed $\begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$, $\begin{pmatrix} -1 & 1 \\ 1 & -1 \end{pmatrix}$, fixed-free $\begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix}$, $\begin{pmatrix} -1 & 1 \\ -1 & 1 \end{pmatrix}$, free-free $\begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix}$, $\begin{pmatrix} -1 & 1 \\ -1 & 1 \end{pmatrix}$.

K 's properties: $\begin{pmatrix} 2 & -1 \\ -1 & 2 & -1 \\ -1 & 2 & -1 \end{pmatrix}$ circular $\begin{pmatrix} 1 & -1 \\ 1 & -1 \end{pmatrix}$, $\begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$, $\begin{pmatrix} 2 & -1 \\ -1 & 2 & -1 \end{pmatrix}$ Δ sim. semipos. $a \lambda = d = 0$, $x_r = (1, 1, 1)$, thus f have $f_1 + f_2 + f_3 = 0$

$\textcircled{1}$ sym. $\textcircled{2}$ tridiag. $\textcircled{3}$ pos. def. ($C_i > 0$ and A has ind. cols) $\textcircled{4}$ K^{-1} is full (not sparse) with all entries > 0 . also $x_r = (1, 1, 1)$ in nullspace.

◦ some proof of Markov's properties:

A 's cols adds to 1 $\rightarrow (1, 1, \dots)^T$ an A^T 's eigenvector with $\lambda = 1$, also A 's $\lambda \rightarrow$ other eigenvalues will be perpendicular to A^T 's $(1, 1, \dots)^T$, which means \rightarrow we need to prove that all other $|\lambda| < 1$.



◻ Perron-Frobenius Theorem:

① $\rho(A) < 1$

• spectrum radius $\rho(A) = \max |\lambda(A)|$, here are some equal conditions: ② $\lim_{k \rightarrow \infty} A^k = 0$

(≈ 2 -norm) (Gelfand Formula: $\rho(A) = \lim_{k \rightarrow \infty} \|A^k\|^{1/k} \leq \|A\|$) ③ $\sum_{k=0}^{\infty} A^k$ converges. (= $(I-A)^{-1}$)

• ($A \geq 0$) spec's monotonicity: $B \in \mathbb{C}$ s.t. $|b_{ij}| < a_{ij} \forall i, j$. then $\rho(B) < \rho(A) < \rho(B)$. Neumann Series with $\rho(A) < 1$ it can converge.

spec with sums: $\sum_{i \rightarrow j} a_{ij} \leq \rho(A) \leq \max_j \sum_i a_{ij}$ Collatz-Wielandt Formula: with $\rho(A) < 1$ it can converge.

• Perron-Frobenius: $A \geq 0$ and irreducible ($1 \leq i \leq n$) $\min \frac{(Ax)_i}{x_i} \leq \rho(A) \leq \max \frac{(Ax)_i}{x_i}$ ('when x is eivec')

then 1. $\rho(A) > 0$ and has $Cu = Au = 1$. 2. all nonneg. eigenvector correspond to eigenvalue $\rho(A)$. (all other λ 's vec has entries < 0)

3. \exists only one $v \in \mathbb{R}^n$ s.t. $v > 0$, $\sum_i v_i = 1$ and $Av = \rho(A)v$; (left eivec) $\dots v^T w = 1$ and $w^T A = \rho(A)w^T$

4. Perron Proj.: $\lim_{k \rightarrow \infty} \left(\frac{A}{\rho(A)}\right)^k = vv^T$ (reducible: $\exists P$ s.t. $P^T A P = \begin{pmatrix} B & C \\ 0 & D \end{pmatrix}$, (or $A^T w = \rho(A)w$) (blockized upper tri.) then A 's graph \times strongly connected.)

◦ sensitivity of eigenvalues: A makes small change ΔA , which results in $\Delta \lambda$:

$$\Delta \lambda = (X + \Delta X)^{-1} (A + \Delta A) (X + \Delta X) - X^{-1} A X = (X^{-1} - X^{-1} \Delta X X^{-1}) A (X + \Delta X) + X^{-1} \Delta A X - X^{-1} A X$$

we use $X^{-1} A = \Lambda X^{-1} \rightarrow = \Lambda X^{-1} \Delta X - X^{-1} \Delta X \Lambda + X^{-1} \Delta A X$. (omit the second term)

now we'll show that $\Lambda X^{-1} \Delta X$ and $X^{-1} \Delta X \Lambda$ cancel on the diag: $\Lambda \subset X^{-1} \Delta X$ with $(X^{-1} \Delta X) \Lambda$

so $\Delta \lambda = X^{-1} \Delta A X$,

◦ proof of $(I-A)^{-1}$: $\rho(A) < 1$ vice above; $\rho(A) = 1$ means a or $\Delta \lambda = y^T \Delta A x$.

$\lambda = 0$, singular; $\rho(A) > 1$ means a $\lambda = \frac{1}{1-\rho} < 0$. since its eigenvector $u > 0$, ($Ax = \lambda x, A^T y = \lambda y$)

◦ matrix series: $\sum_i C_i B^i$ has convergence radius R , then shows there must be entries $(I-A)^{-1} u = \lambda u < 0$

($A \in \mathbb{C}^{n \times n}$) $\sum_i C_i A^i$: ① $\rho(A) < R$, absolute converge; ② $\rho(A) > R$, diverge.

10.3 Markov matrices, population, and economics

Markov mt. : every entry of A is positive $a_{ij} > 0$ ($A > 0$); every cols of A adds to 1.

Then $\lambda_1 = 1$ is larger than any other values, its vec x_1 is the steady state: $u_k = x_1 + \sum_{i=2}^n c_i (\lambda_i)^k x_i$
 $u_{\infty} \rightarrow x_1$.

• Perron-Frobenius Theorem ($\rho(A) = 1$): for $A > 0$, All numbers in $Ax = \lambda_{\max} x$ are positive.

Leslie mt. : population: $\begin{pmatrix} F_1 & F_2 & F_3 \\ P_1 & & \end{pmatrix}$ with $\lambda_{\max} > 1$ maybe. F 's and P 's change in ΔA makes $\Delta \lambda$ (cr/f)

Consumption mt. : $p - Ap = y$, $p = (I - A)^{-1} y$ Neumann/geometric Series:
 (Leontief) product. demand. $(I - A)^{-1} = I + A + A^2 + \dots$

it converges if $\rho(A) < 1$; $\rho(A) = 1$ makes $(I - A)^{-1}$ fail to exist; $\rho(A) > 1$ then $(I - A)^{-1}$ has neg. (cannot meet entries. demand)

10.4 linear programming (A has $m < n$ usually.)

Primal problem : minimize $c^T x$ with $Ax = b$ ($Ax \geq b$) and $x \geq 0$

Dual problem : maximize $y^T b$ with $A^T y \leq c$ and $y \geq 0$

$x \geq 0$ and the slack $s = c - A^T y \geq 0$ gave $x^T s = 0$ which means $y^T b \leq x^T c$, when $y^T b$ max, equality means $x^T s = 0$ thus either $x_j = 0$ or $s_j = 0$, i.e. y solves m eqs. $A^T y = c$ in $c^T x$ min.

geometry: a feasible set ('triangle'), maybe unbounded. the m components that are nonzero in x .

The optimal solution will be one of corners. (m nonzeros and $n - m$ zeros)

① list them and compare. ② the simplex method: start from one corner, enter one variable $0 \rightarrow 1$, others have to adjust to keep $Ax = b$.

③ the interior point method: compute $c \cdot x$ and choose that gives the most neg. change.

remove the constraints $x_j \geq 0$ by then decide how much it can enter and drop one var. to 0 .

minimize $c^T x - \theta \sum \log x_i$ with $Ax = b$. repeat until all changes are pos. (move to neighbor)

Lagrange $L(x, y, \theta) = c^T x - \theta \sum \log x_i - y^T (Ax - b)$

$\frac{\partial L}{\partial x_j} = c_j - \frac{\theta}{x_j} - (A^T y)_j = 0$ which is $x_j s_j = \theta \rightarrow 0$. actually solved by Newton's method (iteration)

10.5 Fourier series: linear algebra for functions

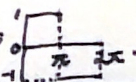
Hilbert space : if and only if their length are finite: $\|v\|^2 = v_1^2 + v_2^2 + v_3^2 + \dots$ (Schwarz ineq. $\sqrt{|v \cdot w|} \leq \|v\| \|w\|$)
 $\|f\|^2 = \int_0^{2\pi} f(x)^2 dx$ — $(f, g) = \int_0^{2\pi} f(x)g(x) dx$

$f(x) = a_0 + a_1 \cos x + b_1 \sin x + a_2 \cos 2x + b_2 \sin 2x + \dots$ (orthon. means $\frac{1}{\sqrt{2\pi}}, \frac{\cos x}{\sqrt{\pi}}, \frac{\sin x}{\sqrt{\pi}}, \dots$)
 basis.

$$\|f\|^2 = 2\pi a_0^2 + \pi(a_1^2 + b_1^2 + a_2^2 + b_2^2 + \dots)$$

since basis are ortho, $a_n = \frac{1}{\sqrt{\pi}} \int_0^{2\pi} f(x) \cos nx dx$, $b_n = \frac{1}{\sqrt{\pi}} \int_0^{2\pi} f(x) \sin nx dx$ ($a_0 = \frac{1}{2\pi} \int_0^{2\pi} f(x) dx = \overline{f(x)}$)

(orthon.)

ex. $\delta(x)$, 

10.6 computer graphics

we use homogeneous coordinates. note to separate 'point' with 'vector'

$$(x, y, z, 1) \quad (x, y, z)$$

Here are some operations: translation $T = \begin{pmatrix} 1 & 0 & 0 & t_x \\ 0 & 1 & 0 & t_y \\ 0 & 0 & 1 & t_z \\ 0 & 0 & 0 & 1 \end{pmatrix}$ (like to rescale the space by $\frac{1}{s}$)
 rescaling $S = \begin{pmatrix} s & 0 & 0 & 0 \\ 0 & s & 0 & 0 \\ 0 & 0 & s & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$; rotation $R = \begin{pmatrix} \cos \theta & 0 & 0 \\ 0 & \cos \theta & 0 \\ 0 & 0 & 1 \end{pmatrix}$ (add two points like to find middle point.)
 (any affine matrix $A = \begin{pmatrix} T(c_1, 0, 0) & 0 \\ T(c_2, 1, 0) & 0 \\ T(c_3, 0, 1) & 0 \\ T(c_4, 0, 0) & 1 \end{pmatrix}$) projection $\begin{pmatrix} I - uu^T & 0 \\ 0 & 1 \end{pmatrix}$ (we use row vec!)
 (mirror $I - 2uu^T$) sometimes there are more steps: ex. proj. to a flat.

10.7 linear algebra for cryptography

modular arithmetic: Inversion of every y ($0 < y < p$) will be possible if and only if p is prime.

Hill cipher: $y_1, 2y_1, \dots, py_1$ have diff remainders, there must be 1. (prime.)
 divide message into blocks $\vec{v}_1, \vec{v}_2, \dots$ of size n .
 multiply each by encryption mt. $E \pmod p$. decryption mt. will be E^{-1} (p is prime and this exists.)

finite field $F_p = \{0, 1, 2, \dots, p-1\}$ or p^k members. define its $+$, \times rules. (ex. mod) (It can be difficult when p is large or done many times.)

F_2^3 have 8 vecs. 29 mts. over F_2 (3×3), maybe sin.

11. Numerical Linear Algebra

speed, accuracy, stability

11.1 Gaussian elimination in practice

roundoff error \rightarrow partial pivoting: choose the largest num in row k or below, exchange

for pos. definite mt, row exchanges aren't required. (no improve)

operation counts of full mt.: Gaussian elim. has two advantages over $A^{-1}b$:

orthogonalization: $\frac{2}{3}n^3$ (Householder) $\frac{1}{3}n^3$ mul-sub compared to n^3 in A^{-1} (next n^2 equal)

(Givens) $\textcircled{2}$ if A is banded so are L, U , but A^{-1} is full.

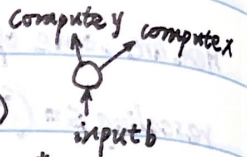
Δ Euler angles: $\alpha_{32}, \alpha_{31}, \alpha_{21} A = R$, where A is ortho so $R = I$; $A = \alpha_{21}^{-1} \alpha_{31}^{-1} \alpha_{32}^{-1}$

□ LA in probability and statistics:

线性在概率论中(简)

Monte Carlo method: $E(x) \leftarrow \frac{\sum x_i}{N}$, $\sigma \sim \frac{1}{\sqrt{N}}$

multilevel MC: 2-level $E(x) = \frac{1}{N} \sum_{i=1}^N y(b_i) + \frac{1}{N^*} \sum_{i=1}^{N^*} (x(b_i) - y(b_i))$



(for same accuracy, since $x-y$ has a smaller σ^* , N^* can smaller than N)

fixed cost $NC + N^*C^* = T$, optimal ratio $\frac{N^*}{N} = \frac{\sigma^*}{\sigma} \sqrt{\frac{C}{C^*}}$

3-level $E(x) = \frac{\sum x}{N} + \frac{\sum (y-z)}{N^*} + \frac{\sum (x-y)}{N^{**}}$... and optimize those N .

Covariance matrix: $V = E((x-\bar{x})(x-\bar{x})^T)$ ($\sum \sum$ or $\int dx$)

V is sym. positive semidefinite (most pos., unless the experiments are dependent)

correlation mt. $u^T V u = E(\|u^T(x-\bar{x})\|^2) \geq 0$. (sim. $|V| = 0$)

$R = \begin{pmatrix} 1 & \rho_{xy} \\ \rho_{yx} & 1 \end{pmatrix} = DVD$ where $D = \text{diag}(\frac{1}{\sigma_i^2})$. $\rho_{xy} = \frac{\sigma_{xy}}{\sigma_x \sigma_y} \in [-1, 1]$ and we have $\sigma_i^2 \sigma_j^2 \geq \sigma_{ij}^2$ transform for $Z = AX$.

(multivariate) Gaussian $p(x) = \frac{1}{(\sqrt{2\pi})^m |V|} e^{-\frac{(x-m)^T V^{-1} (x-m)}{2}}$, $m = \text{means of } x$.

we can diagonalize the cov. mt $V^{-1} = Q \Lambda^{-1} Q^T$, $Y = Q^T X$, $X = Q Y$, which means finding

verify: $\int p(x) dx = \int e^{-\frac{Y^T \Lambda^{-1} Y}{2}} \frac{1}{\sqrt{2\pi}^m} dY = \prod \frac{1}{\sqrt{2\pi} \lambda_i} = \frac{1}{\sqrt{2\pi}^m |V|}$ ($dx = dY$) combs. that are ind. exper.

$$\int x p(x) dx = \int x p(x) dx + m \int p(x) dx = m$$

$$\int (x-m) p(x) (x-m)^T dx = Q \int Y Y^T e^{-\frac{Y^T \Lambda^{-1} Y}{2}} dY Q^T = Q \Lambda Q^T = V$$

weighted least squares:

previously we choose \hat{x} minimize $\|b - Ax\|^2$, get

now we minimize $E = (b - Ax)^T V^{-1} (b - Ax)$

$ATA \hat{x} = ATb$. \leftarrow normal dis. $\sigma^2 = 1, m=0$.

'whitening noise' $ATV^{-1}Ax = ATV^{-1}b$ (if b 's errors are not ind. or variances not equal.)

(means $\frac{b_i}{\sigma_i}$ to get $N(0,1)$. then $V^{-\frac{1}{2}}Ax = V^{-\frac{1}{2}}b$ (min $\sum_{i=1}^m \frac{(b-Ax)_i^2}{\sigma_i^2}$)

substitute by $A \rightarrow V^{-\frac{1}{2}}A$ and $b \rightarrow V^{-\frac{1}{2}}b$)

variance of \hat{x} : since $\hat{x} = Lb = (ATV^{-1}A)^{-1} ATV^{-1}b$ use trans. formula to get $V(\hat{x}) = (ATV^{-1}A)^{-1}$

recursive least squares: static / dynamic: Kalman filter:

update: $\begin{pmatrix} A_0 \\ A_1 \end{pmatrix} \hat{x}_1 = \begin{pmatrix} b_0 \\ b_1 \end{pmatrix}$ (without $x_k = F_{k-1} x_{k-1}$, C_k changing state here)

$$\hat{x}_1 = \hat{x}_0 + K_1 (b_1 - A_1 \hat{x}_0)$$

cov. mt. $W_1^{-1} = W_0^{-1} + A_1^T V_1^{-1} A_1$ (the same form)

Kalman gain mt. $K_1 = W_1 A_1^T V_1^{-1}$

11.2 norms and condition numbers

we choose ℓ^2 norm here. sym. $A: \|A\| = |\lambda|_{\max}(A) (= \rho(A))$.

sensitivity to error: sym. or unsym. $A: \|A\| = \sqrt{\lambda_{\max}(A^T A)} (= \sigma_{\max}(A))$.

$$Ax = b \quad \rightarrow \quad A(x + \Delta x) = b + \Delta b \quad \frac{\|\Delta x\|}{\|x\|} \leq c \frac{\|\Delta b\|}{\|b\|} \quad \text{condition number } c = \|A\| \|A^{-1}\| \geq 1$$

$$\downarrow \quad (A + \Delta A)(x + \Delta x) = b \quad \frac{\|\Delta x\|}{\|x\|} \leq c \frac{\|\Delta A\|}{\|A\|} \quad (\text{for pos. mt, } c = \frac{\lambda_{\max}}{\lambda_{\min}})$$

equal when b along largest eigen, Δb along smallest eigen, amplified by c .

11.3 iterative methods and preconditioners

) split A into $S-T$. iteration $Sx_{k+1} = Tx_k + b$ — two goals: speed per step and fast converg.

1. Jacobi method: keep the diag (A) as S , off-diag part is T . $\rho < 1$ (diag) or not.)

2. Gauss-Seidel method: keep the lower tri. part as S . error eq. $e_{k+1} = S^{-1} T e_k$

& SOR (successive overrelaxation): (cut storage and usually speed up iter.) S has diag(A) and below diag of ωA . ($e_i = x_{\infty} - x_i$)

3. Elimination (exact LU) (triang. very fast.) choose ω to make the $\rho(S^{-1}T)$ small asp.

& incomplete LU: ('fill-in'. sparse mt. that has nonzeros far from diag.)

4. multigrid 5. conjugate gradients & preconditioned CG. $L_0 U_0 x_{k+1} = (L_0 U_0 - A) x_k + b$ set small nonzeros to 0 in L.U.

for eigens: 6. power method & inverse power method. (Arnoldi iter.; Lanczos iter.)

$u_k = A^k u_0$ the largest eivalue dominates. u_k gradually points towards x_1 ;

or by A^{-1} , solve $Au_{k+1} = u_k$ (by saving L.U.), λ_{\min} dominates. For speed we shifted:

8. QR method: $A = QR$, $A_1 = RQ = Q^{-1} A Q$ thus same eivalue with A . $A_1 = Q_1 R_1$, $A_2 = R_1 Q_1$, ... eivalue with A . even smaller. $A - \lambda^* I$ (λ^* can choose a to make λ_{\min} Rayleigh quo. $\frac{x^T A x}{x^T x}$)

two extra ideas:

- ① shift. $A_k - c_k I$ into $Q_k R_k$. $A_{k+1} = R_k Q_k + c_k I$ back. (c_k choose to near an (unknown) eivalue through diag of A_k .)
- ② elimination. obtain zeros before. EAE' (or Givens) (move) leave nonzeros along subdiag, which QR factor. drops from (Hessenberg mt.) $O(n^3) \rightarrow O(n^2)$.